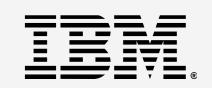
## OpenSSL Conference Prague 2025

# In-Situ Performance Measurements of Crypto-Algorithms in TLS v1.3

Dr. Martin L. Schmatz

IBM Research - Zurich

Oct 09, 2025



#### Motivation

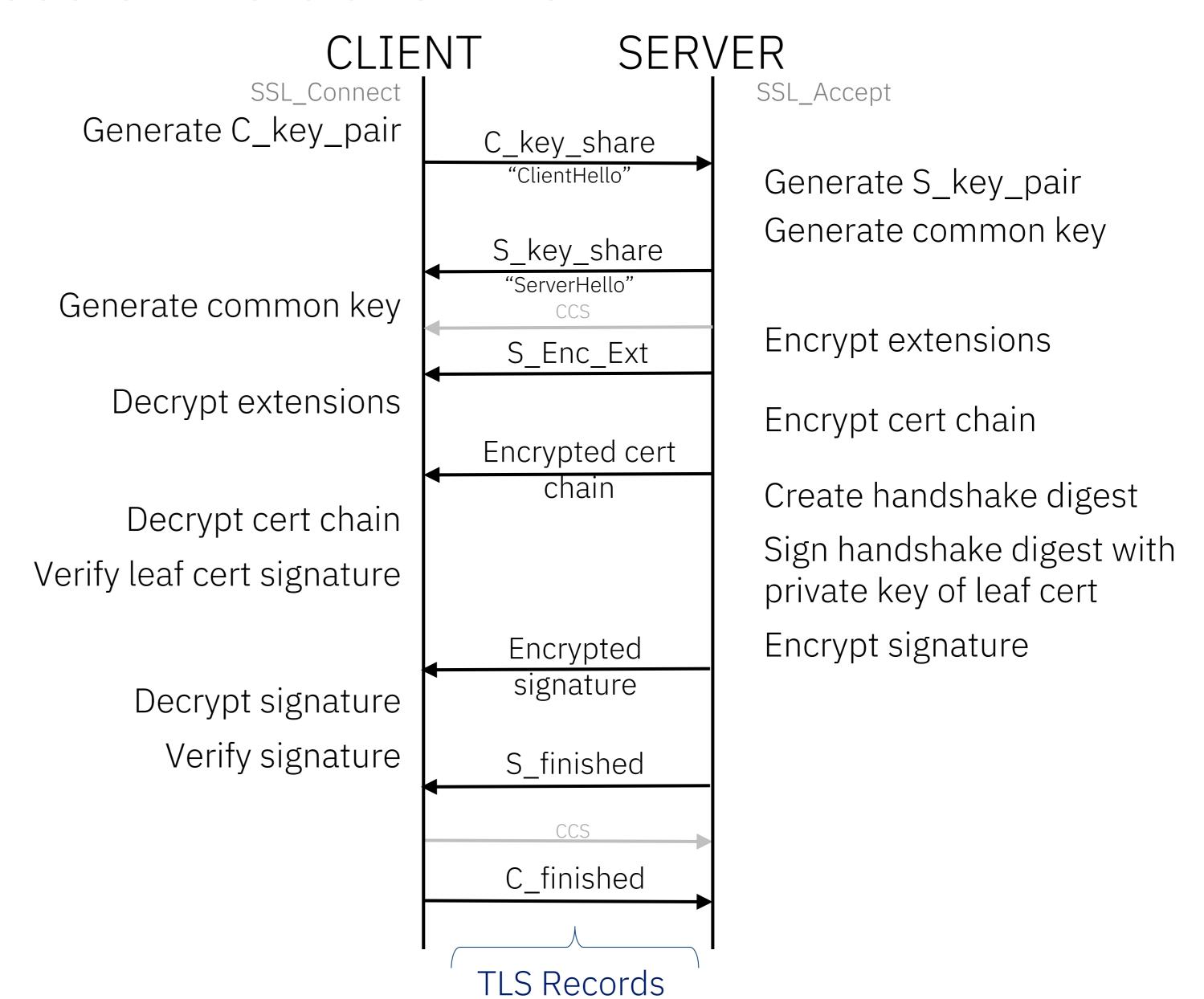
Clients are concerned about the *processing efforts* when using *Q-safe TLS* connections for:

- Key establishment
- Authentication (~= certificates)
- Confidentiality & Authenticity (~=enc-/de-cryption)

Clients are concerned about the required network bandwidth when using *Q-safe TLS* connections

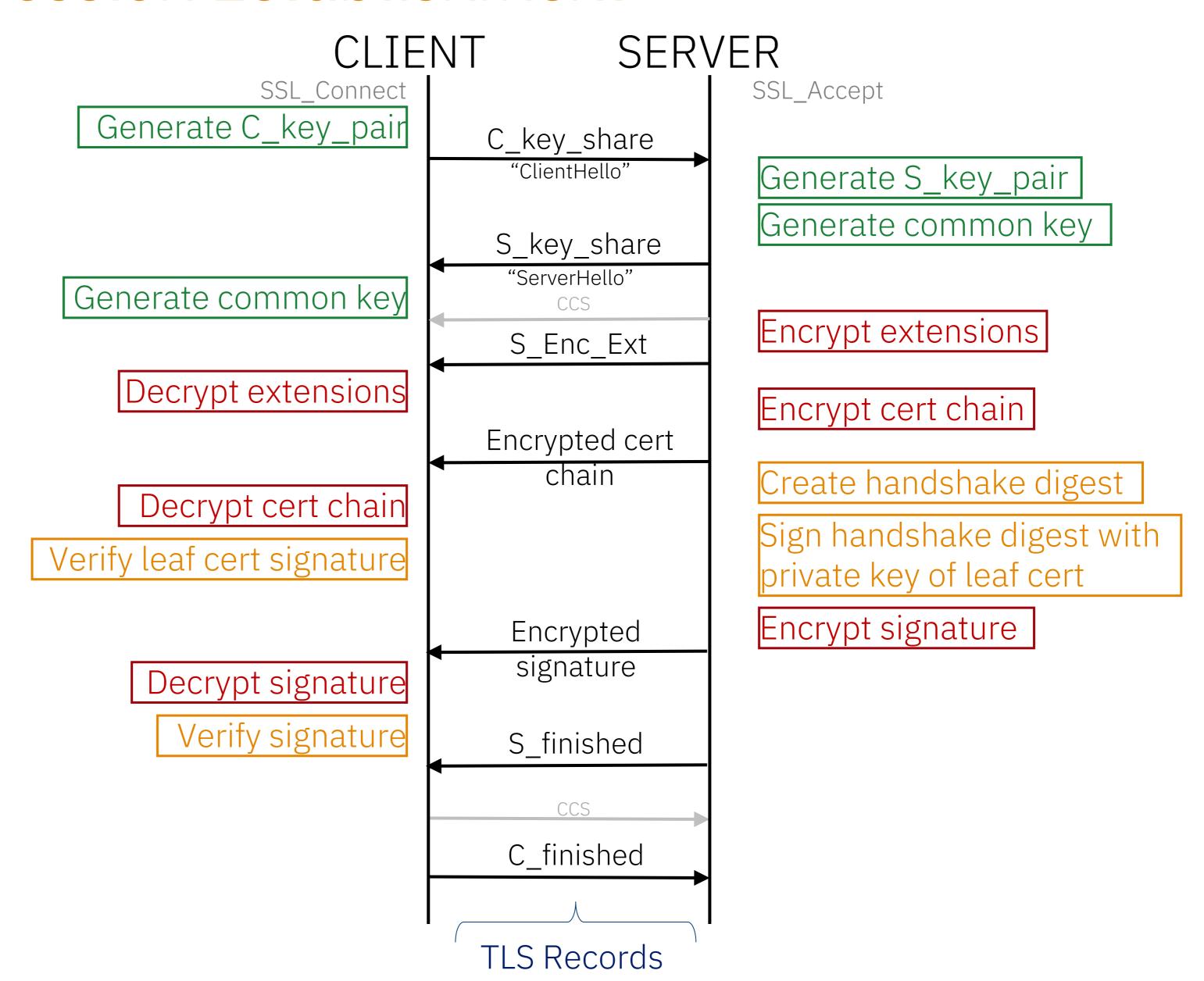
- While synthetic (cycle counting) and measured performance for the processing of Q-safe algorithms are readily available, those DO NOT include collateral performance penalties for the use of the algorithms in the TLS v1.3 protocol into account (e.g., memory allocations, data copying, algorithm loading, merging hybrid secrets, etc.)
- Hence the desire to get *in-situ performance measurements:* <u>CPU time</u> required to execute the individual steps of a TLS protocol session establishment and the data volume to be transported over the network as a function of Q-safe algorithms.

Only relevant parts shown, no HelloRetryRequest (key share fits supported groups)



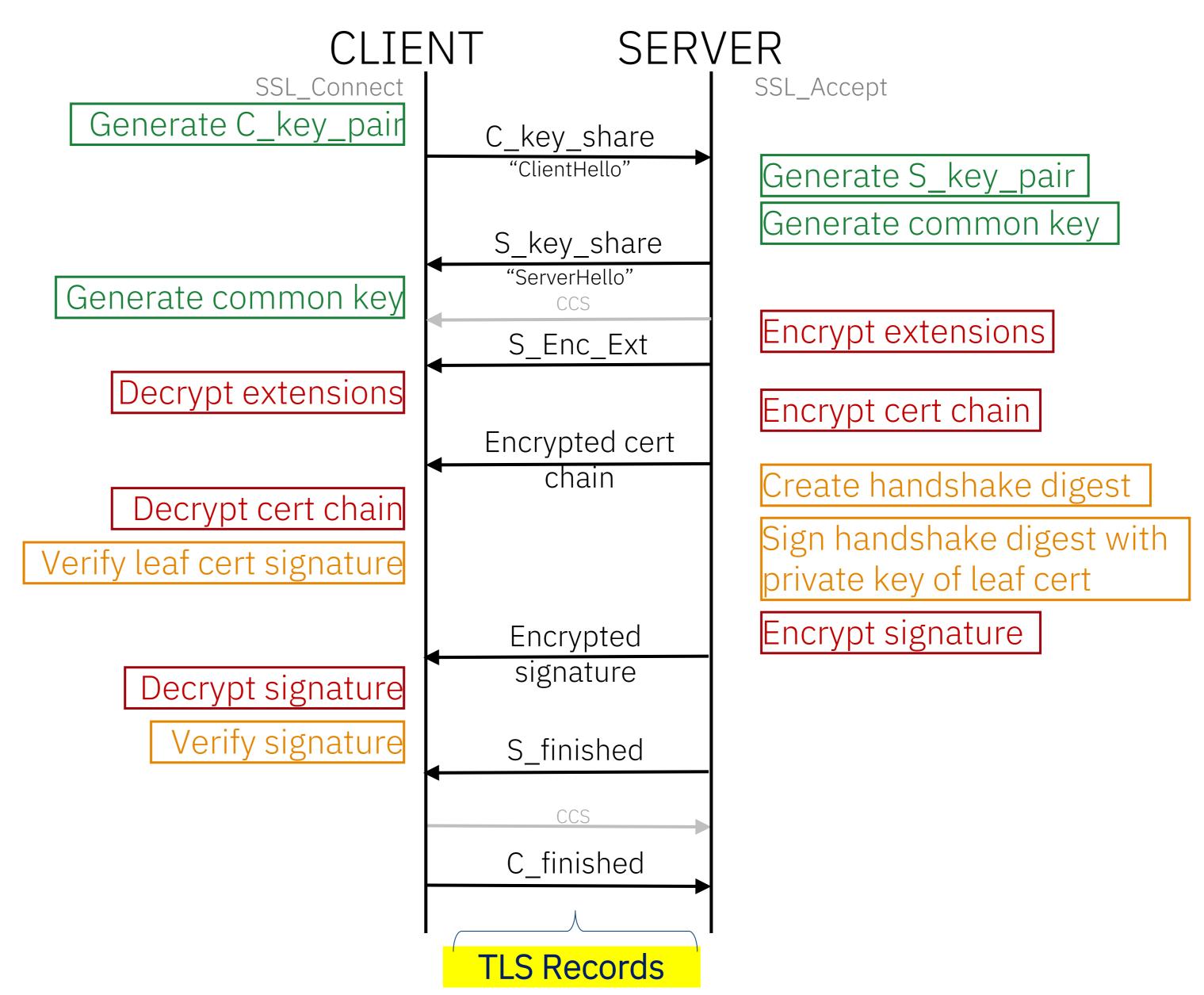
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Some processing is for key establishment, for confidentiality & authenticity, and for authentication



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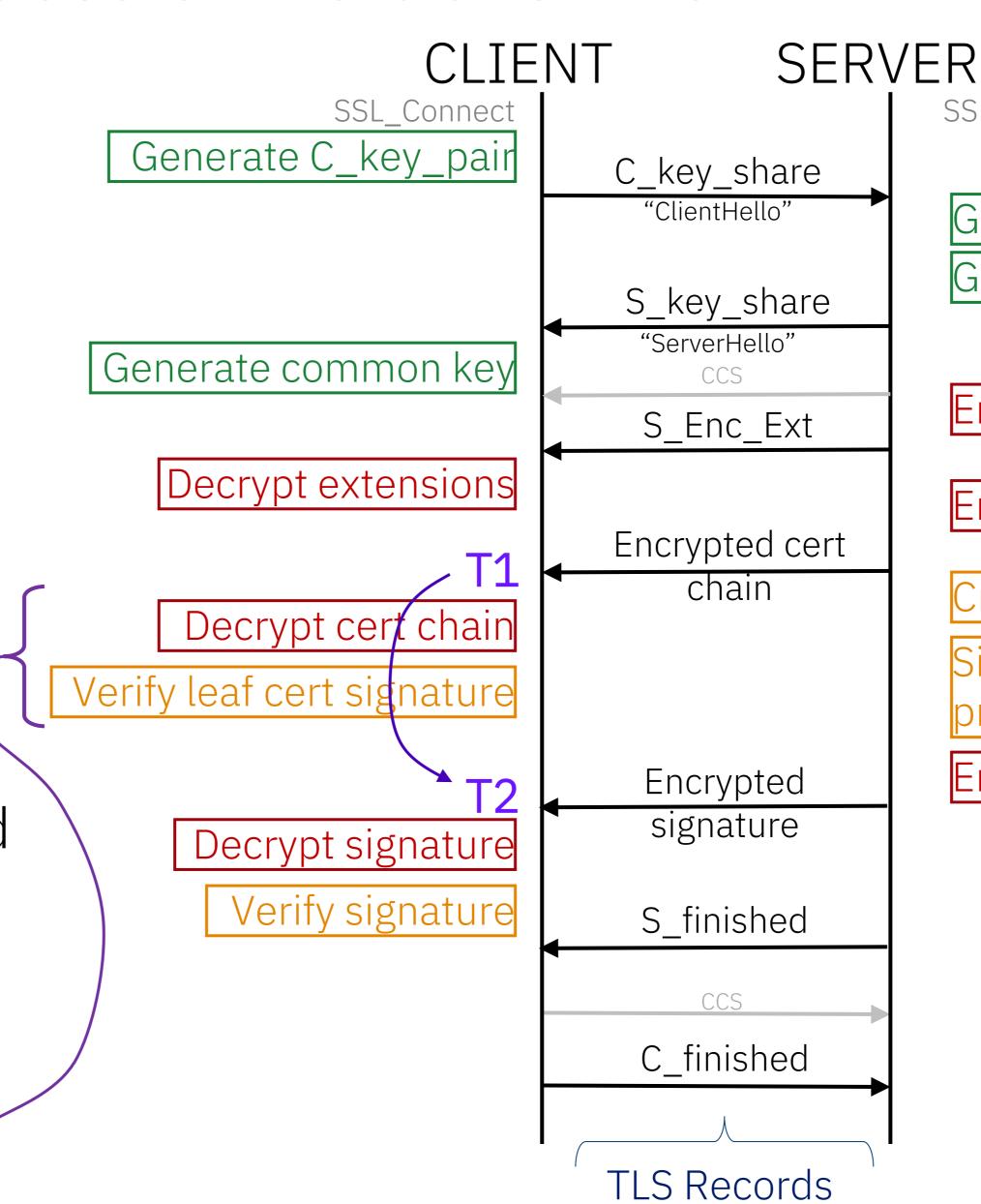
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Some processing is for key establishment, for confidentiality & authenticity, and for authentication

The measured timestamps of TLS record *processing* start/end can be used to get CPU time used for the computations

#### Example:

T2-T1 = Time to decrypt cert chain and to verify cert signature(s)



SSL\_Accept

Generate S\_key\_pair

Generate common key

Encrypt extensions

Encrypt cert chain

Create handshake digest

Sign handshake digest with private key of leaf cert

Encrypt signature

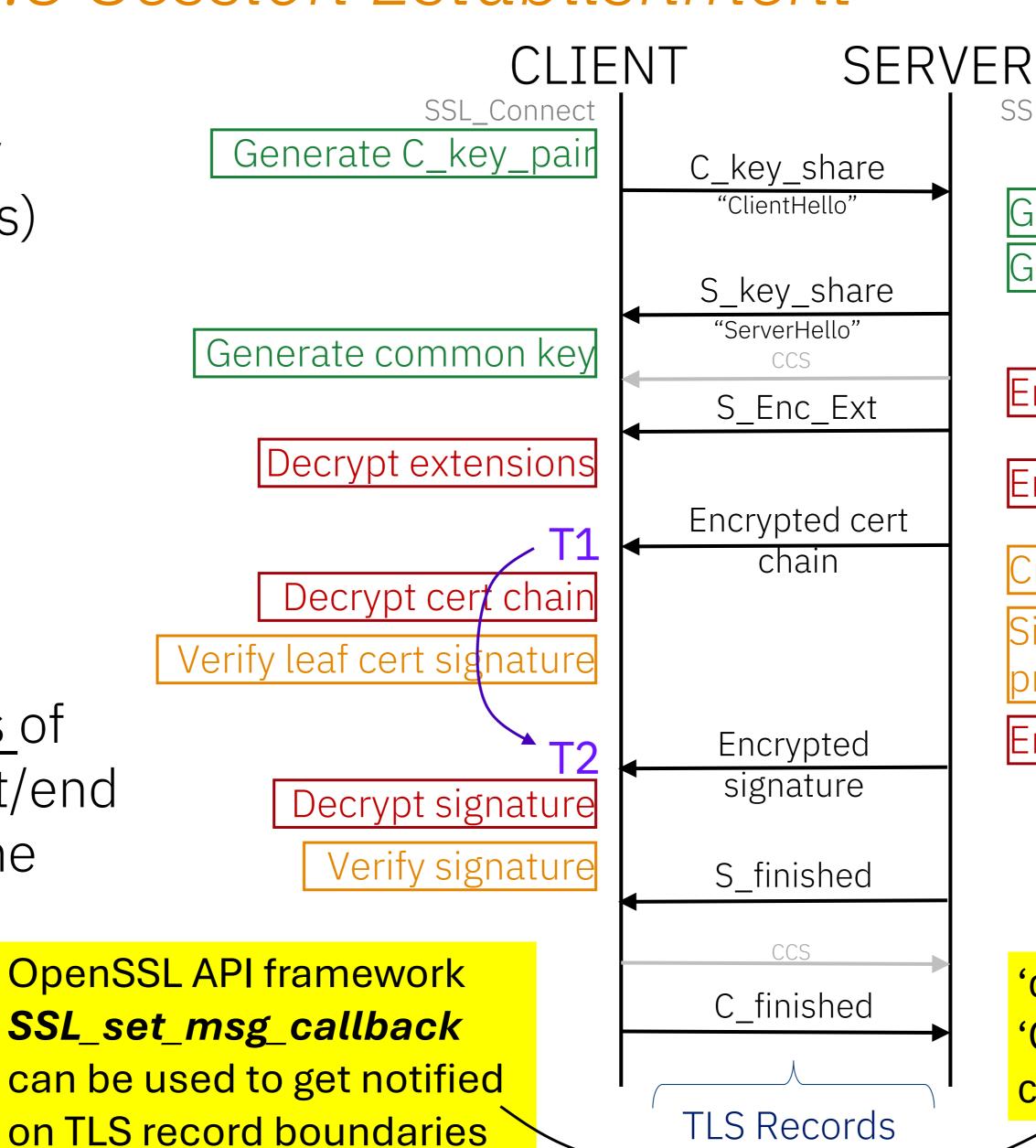
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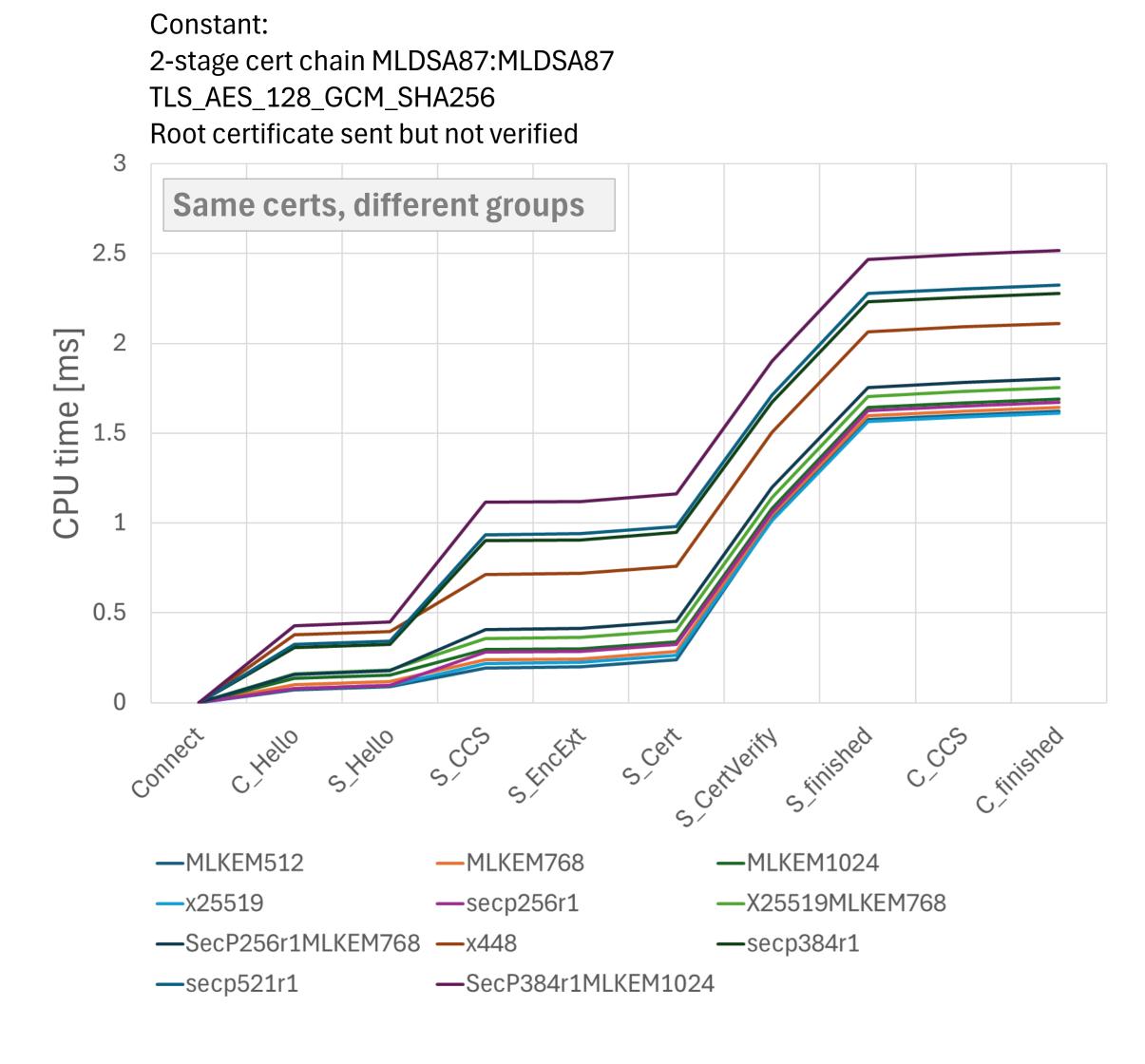
Create handshake digest

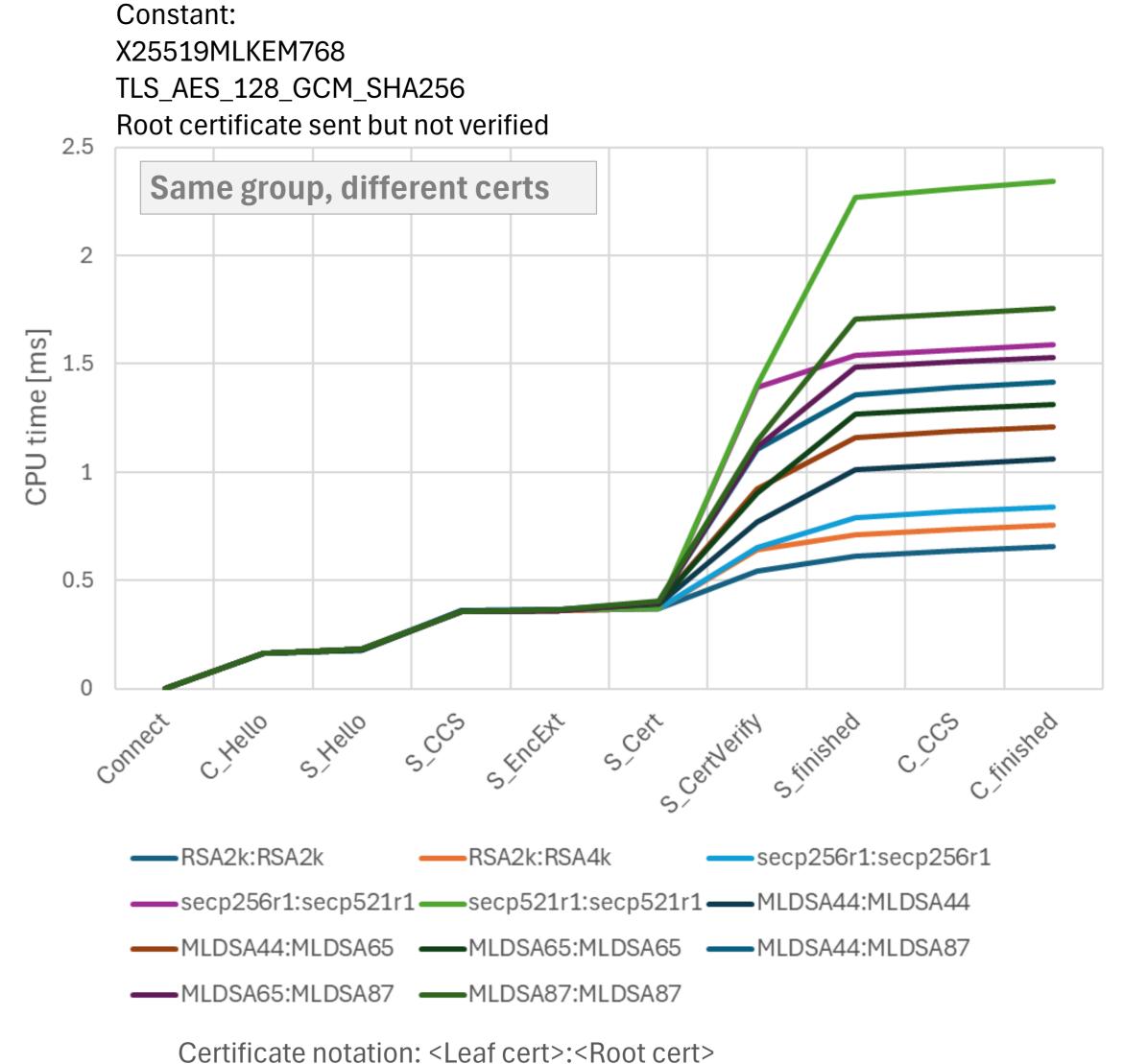
Sign handshake digest with private key of leaf cert

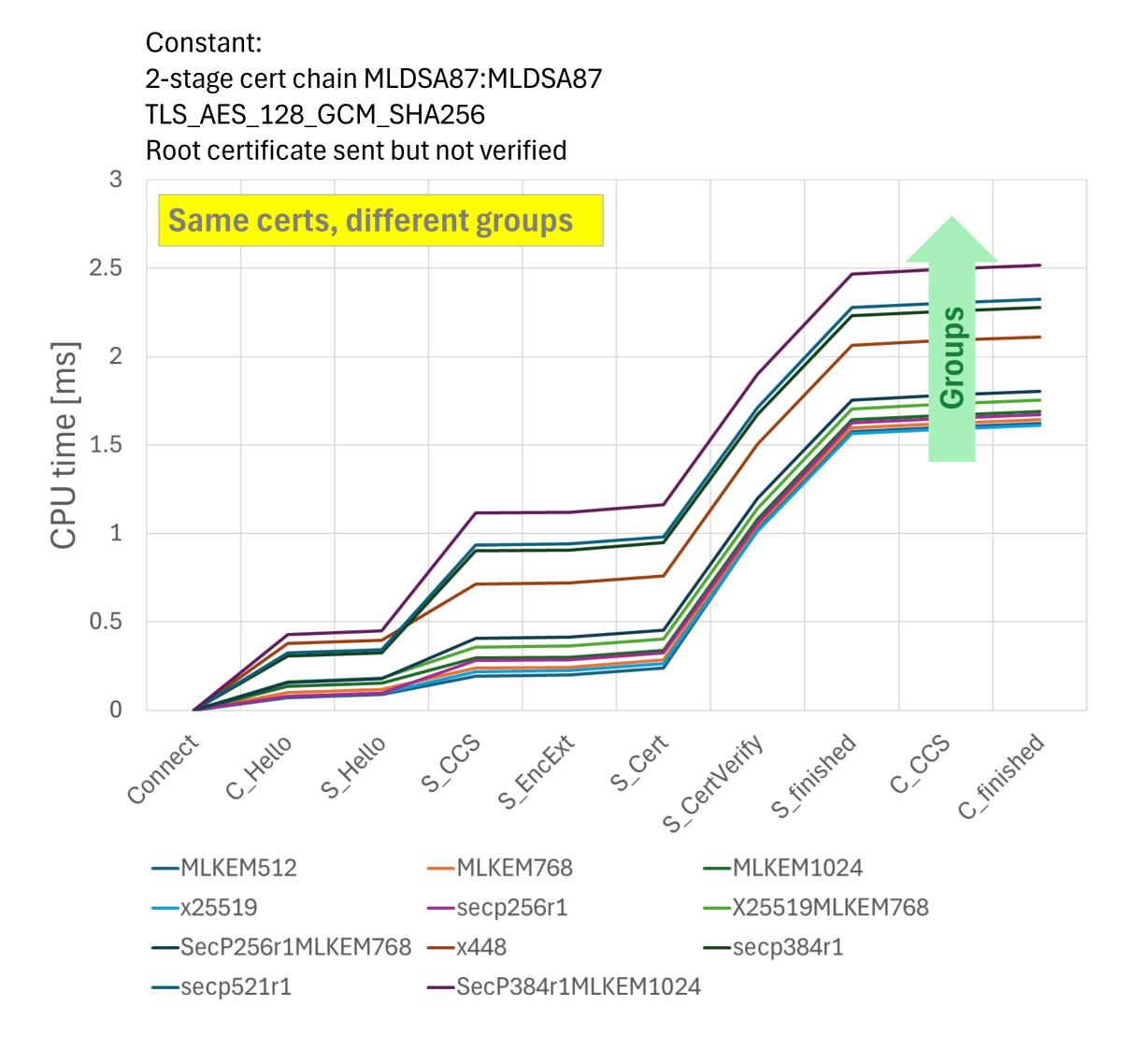
Encrypt signature

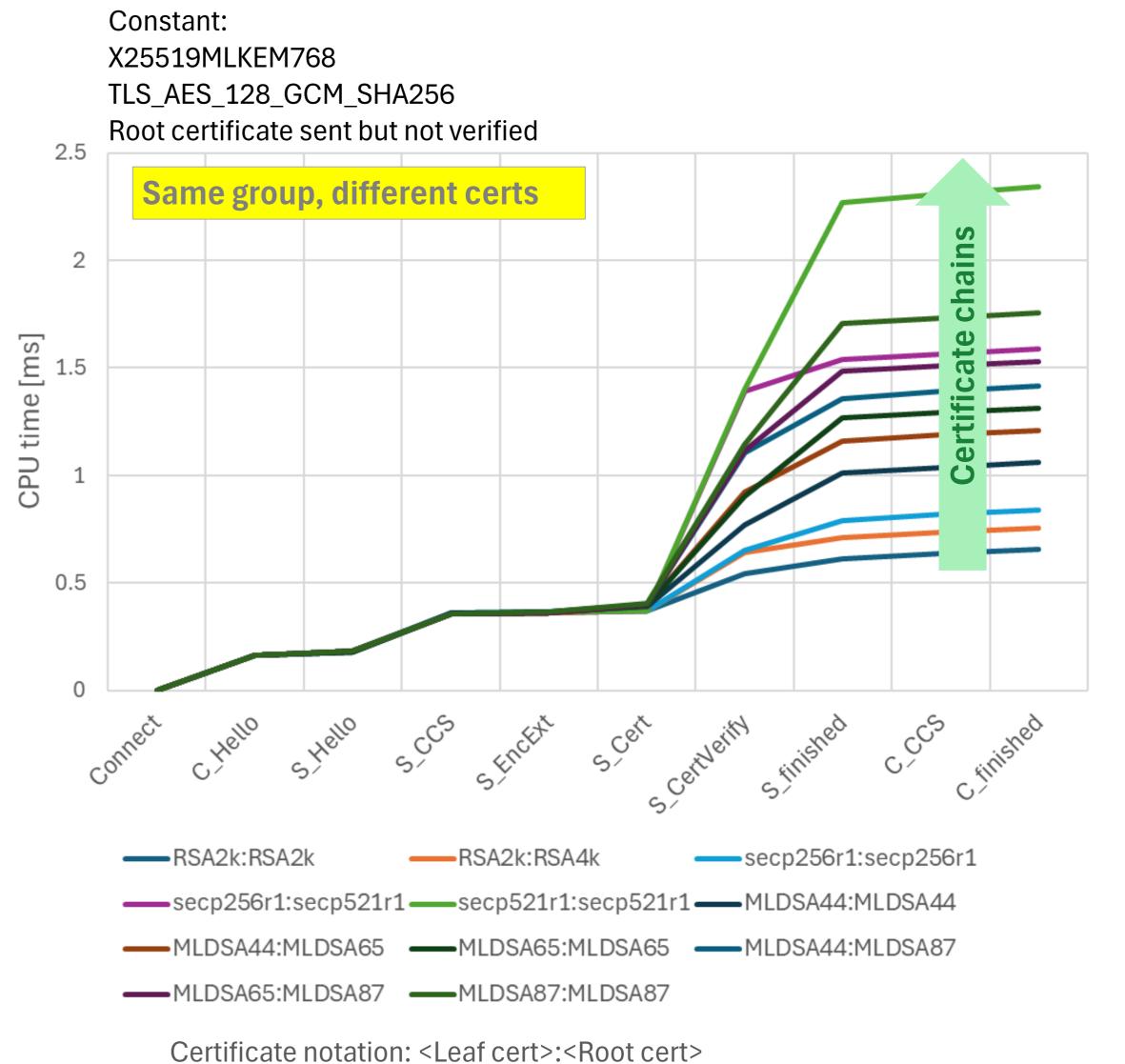
'clock\_gettime' function with 'CLOCK\_THREAD\_CPUTIME\_ID' can get (accumulated) CPU time

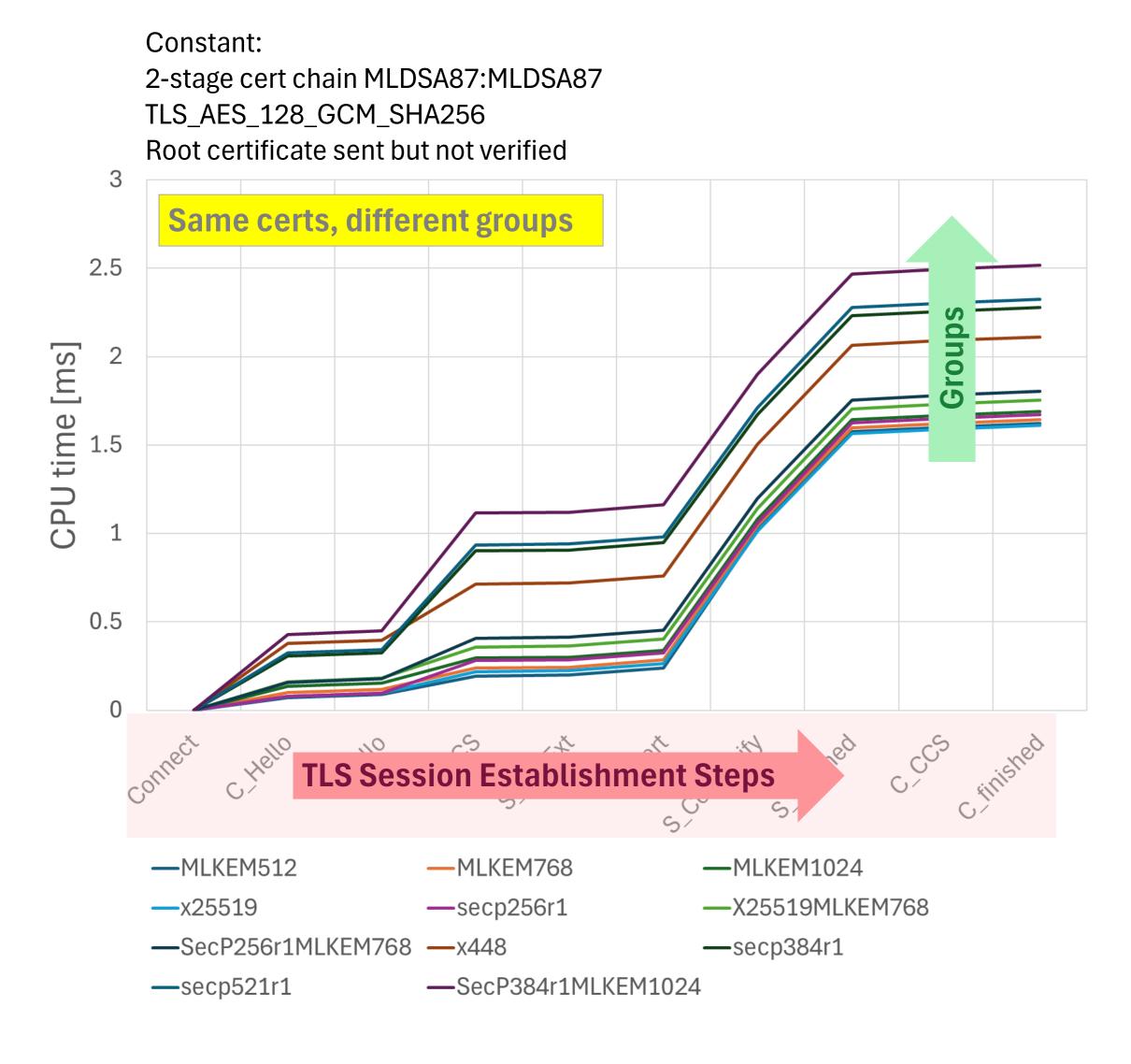
## Measurement examples: <u>CLIENT</u>

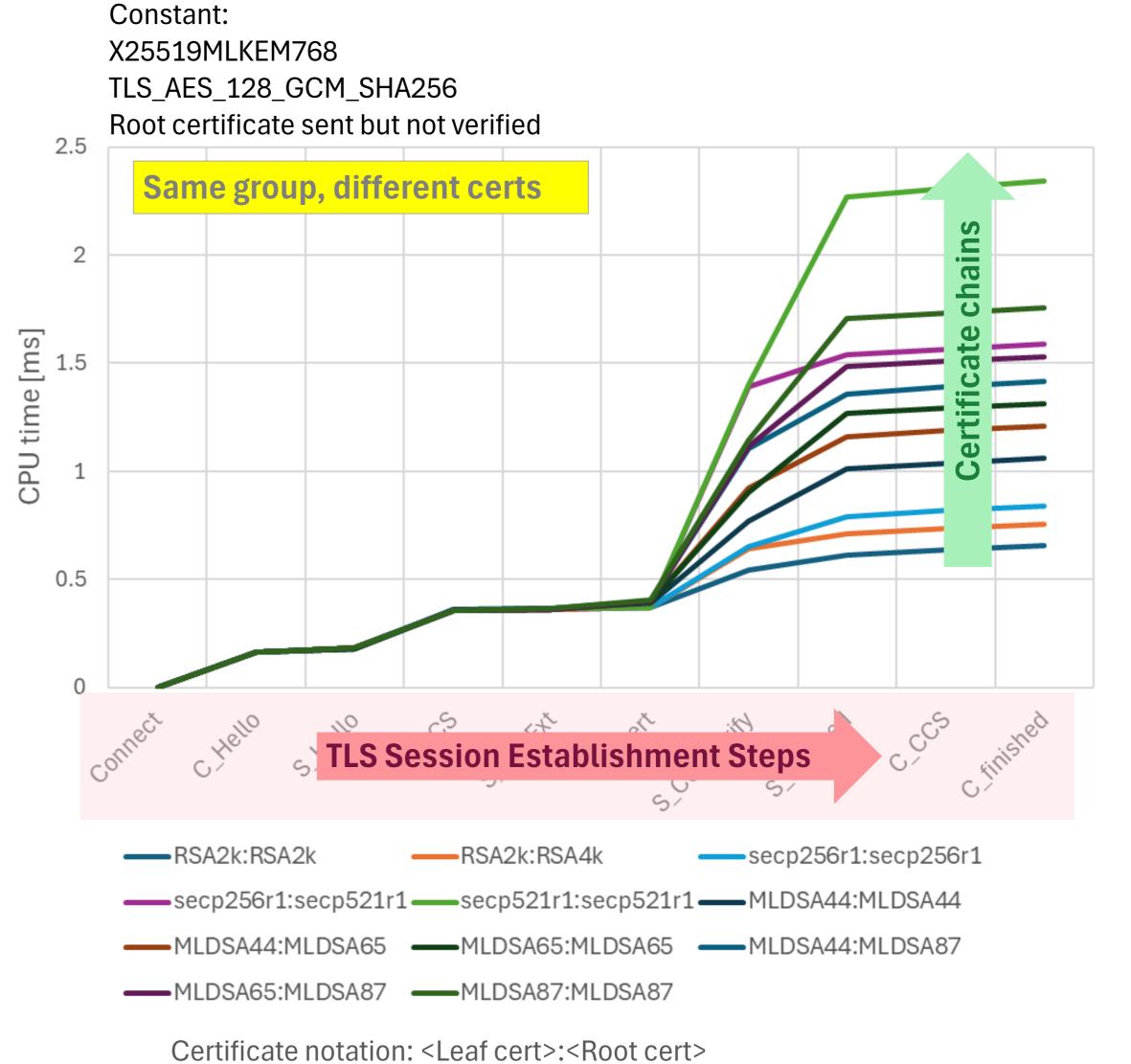


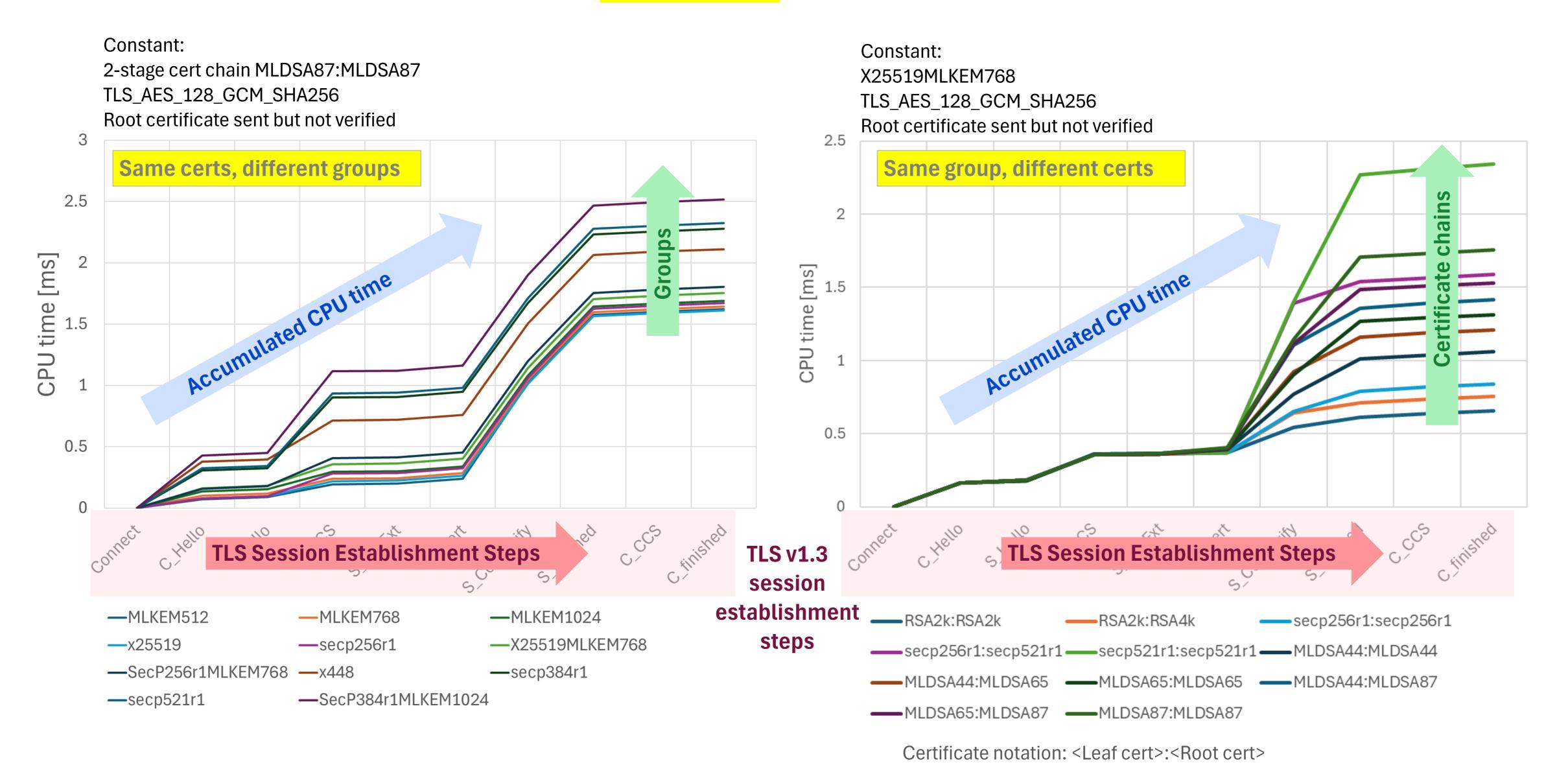


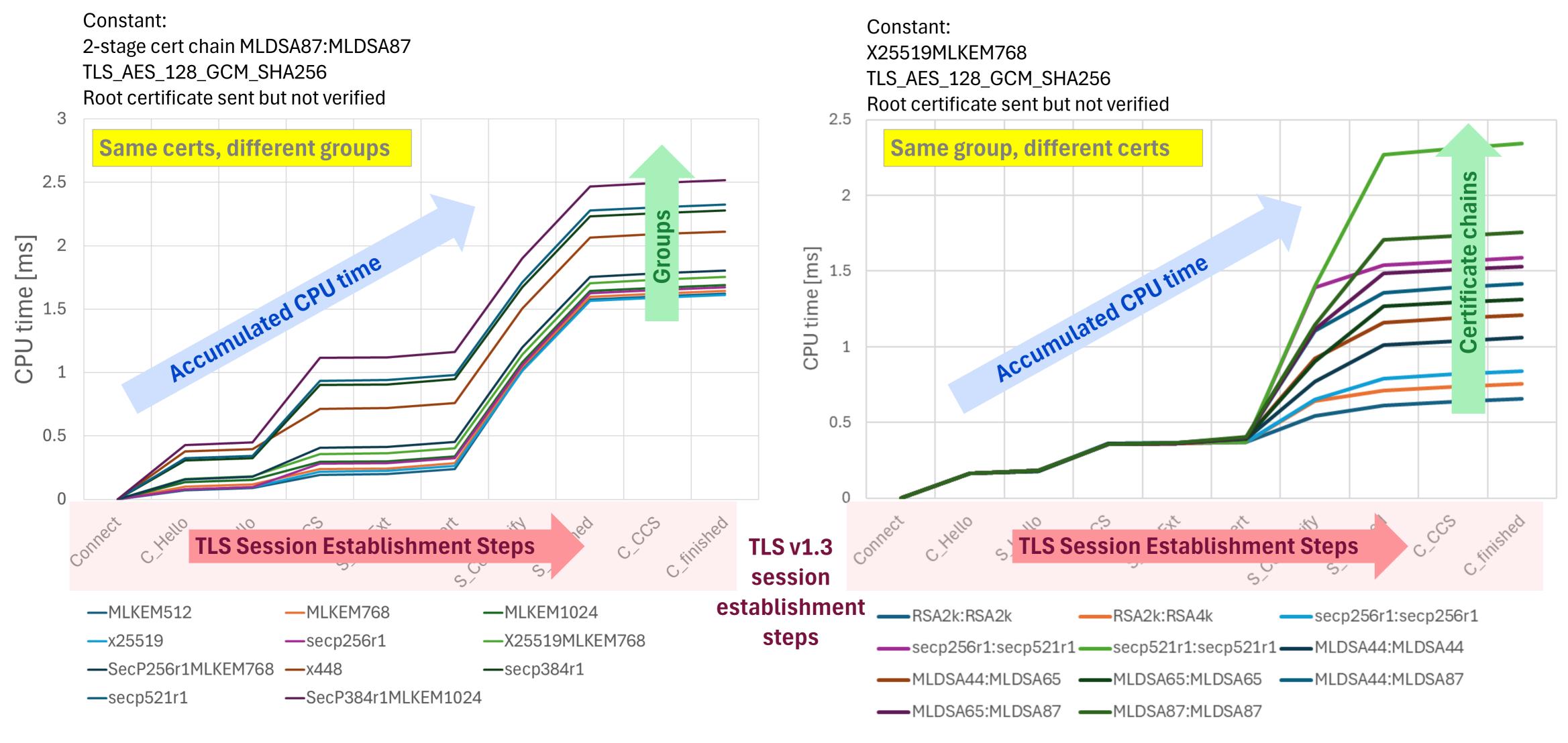






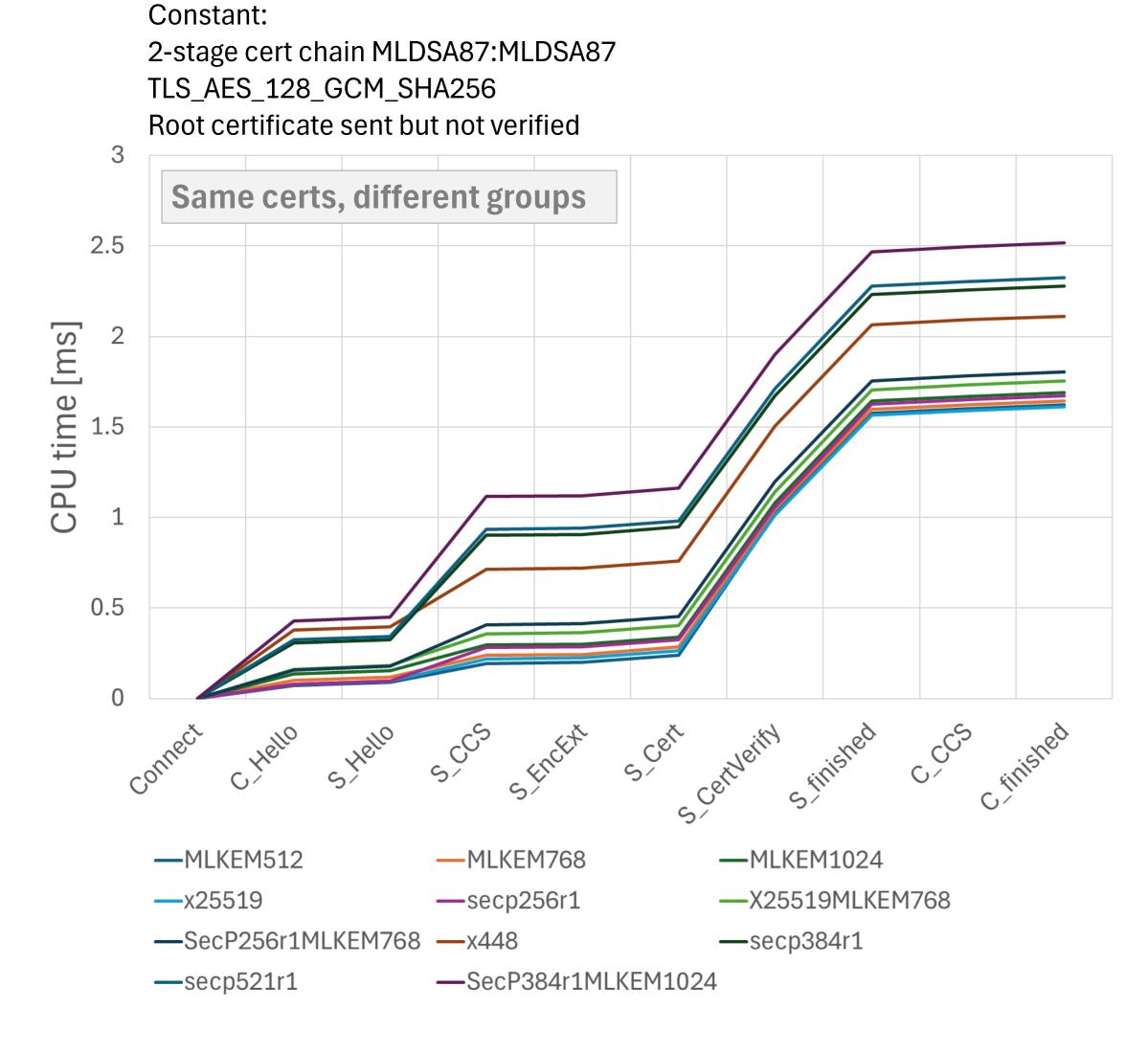


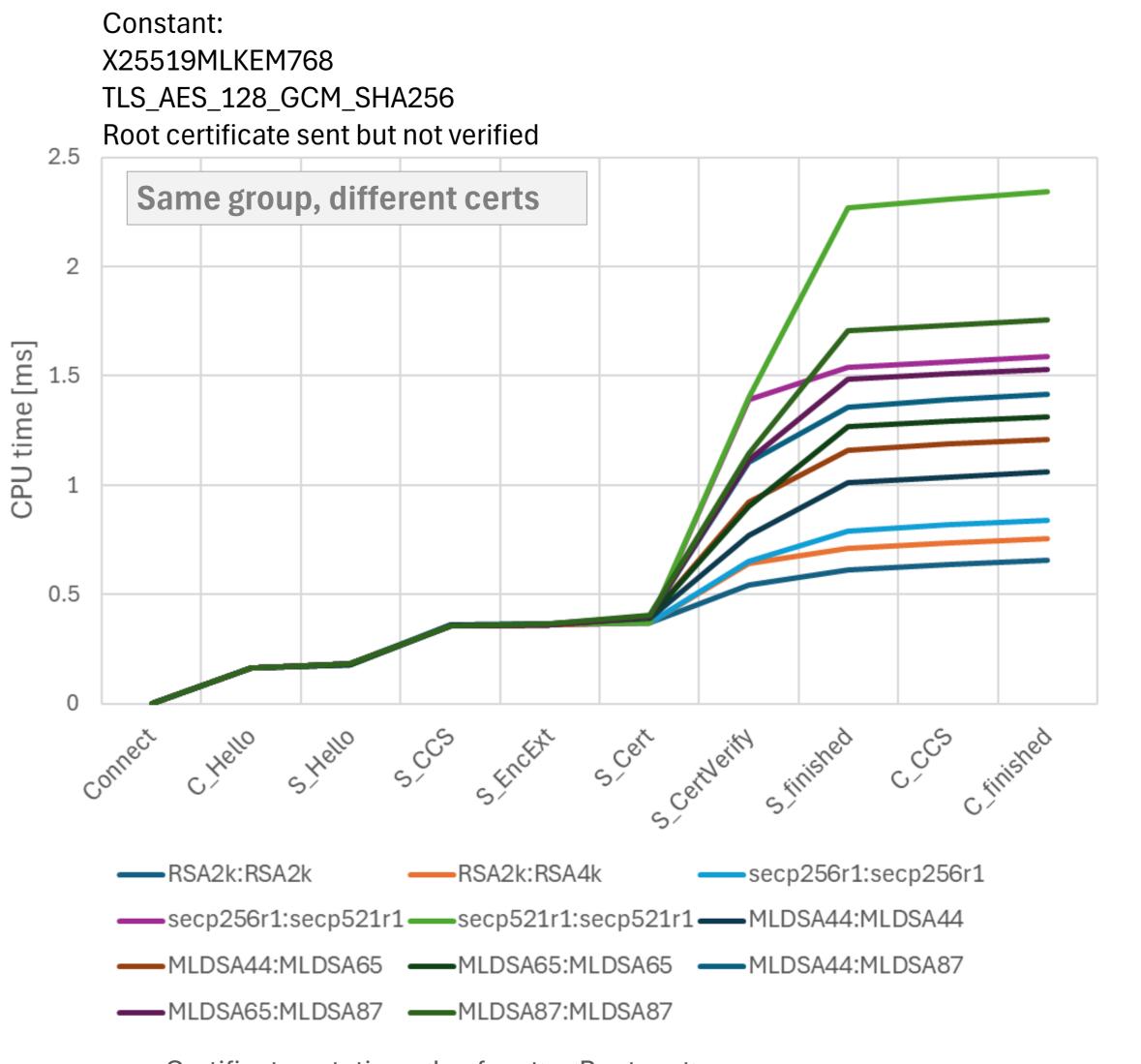




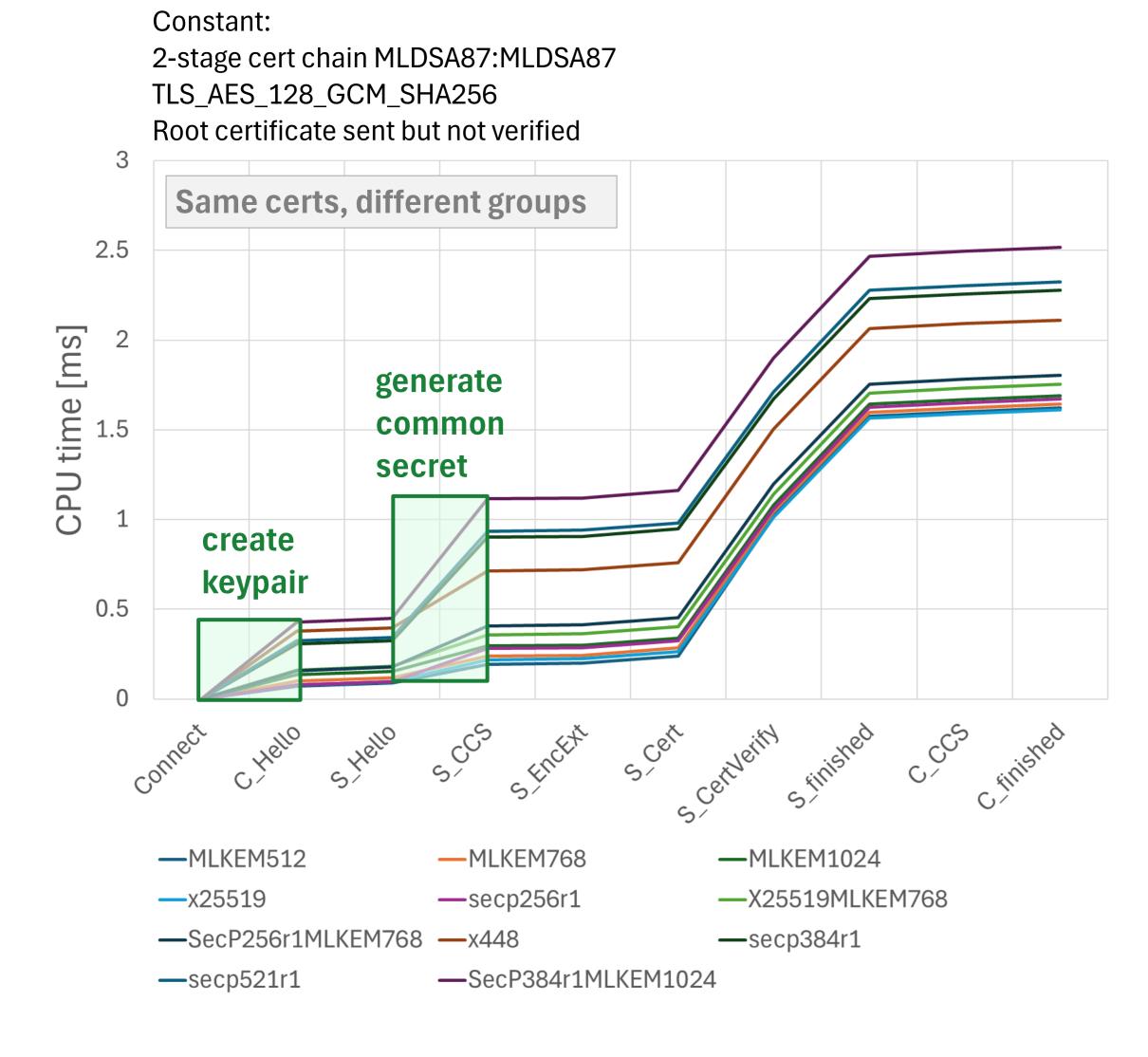
Certificate notation: <Leaf cert>:<Root cert>

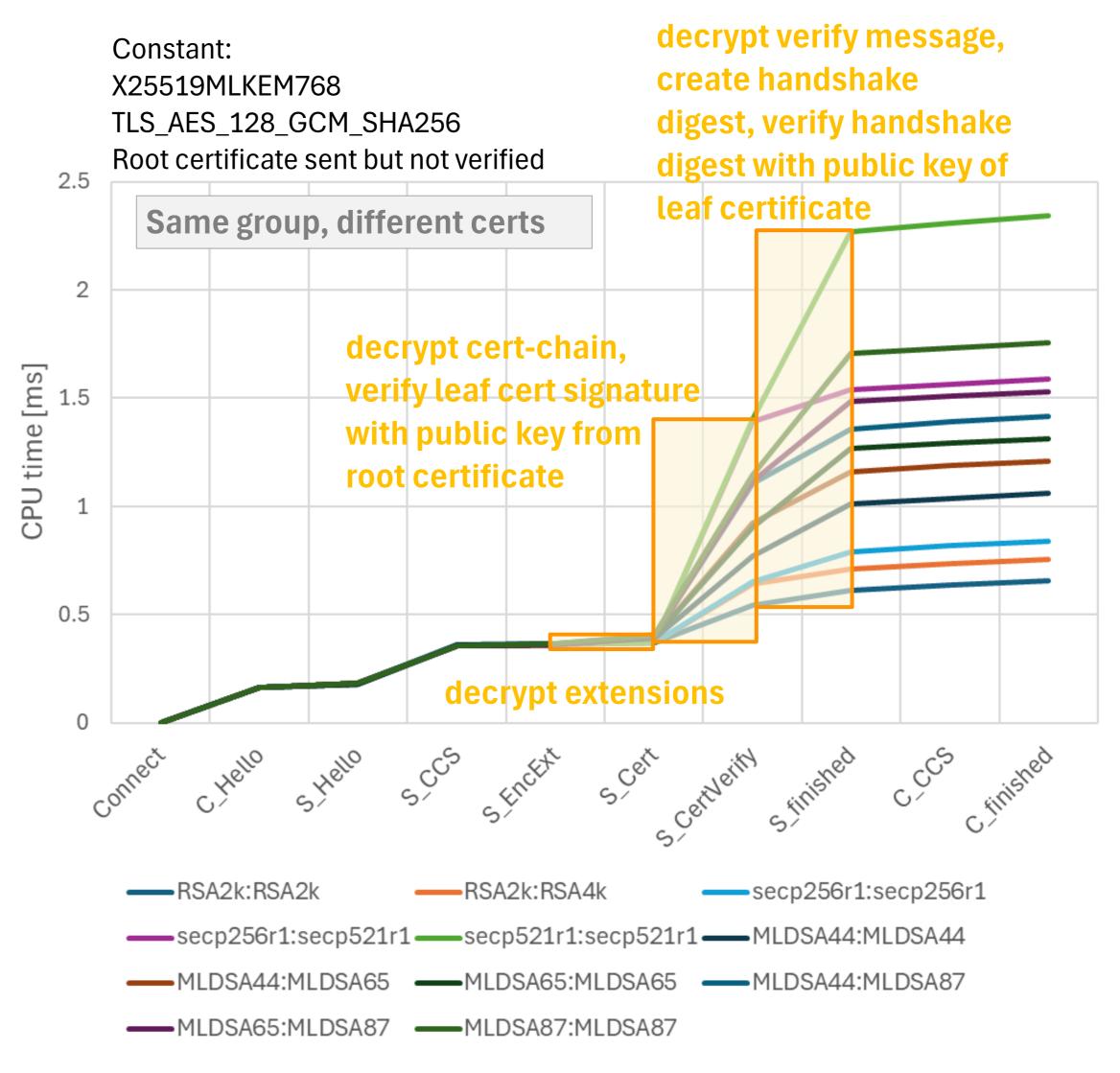
# Data extraction: CLIENT





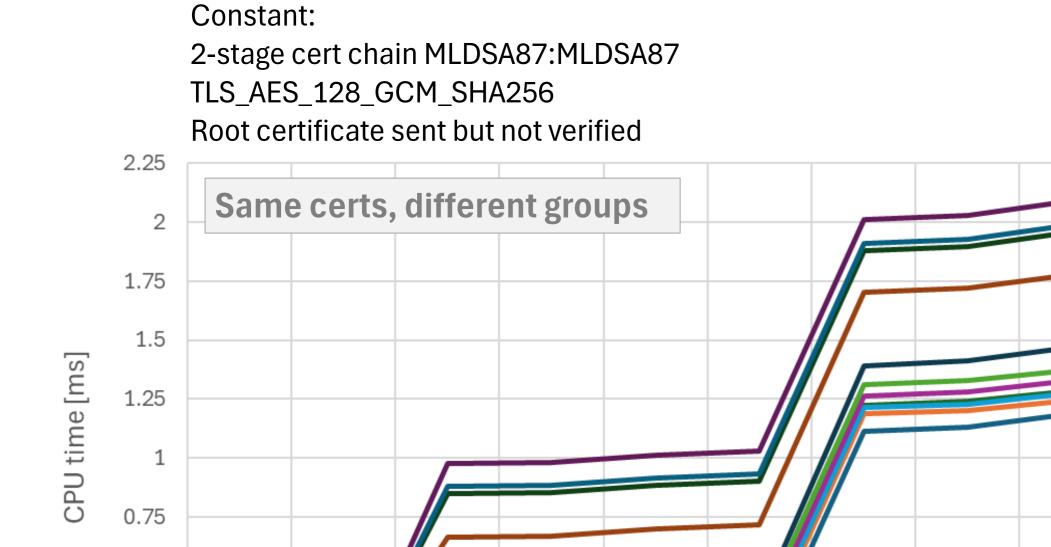
## Data extraction: CLIENT





Certificate notation: <Leaf cert>:<Root cert>

# Data extraction: SERVER



-MLKEM768

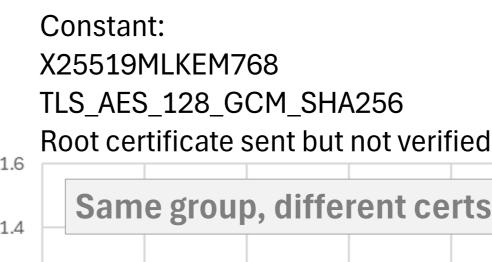
---secp256r1

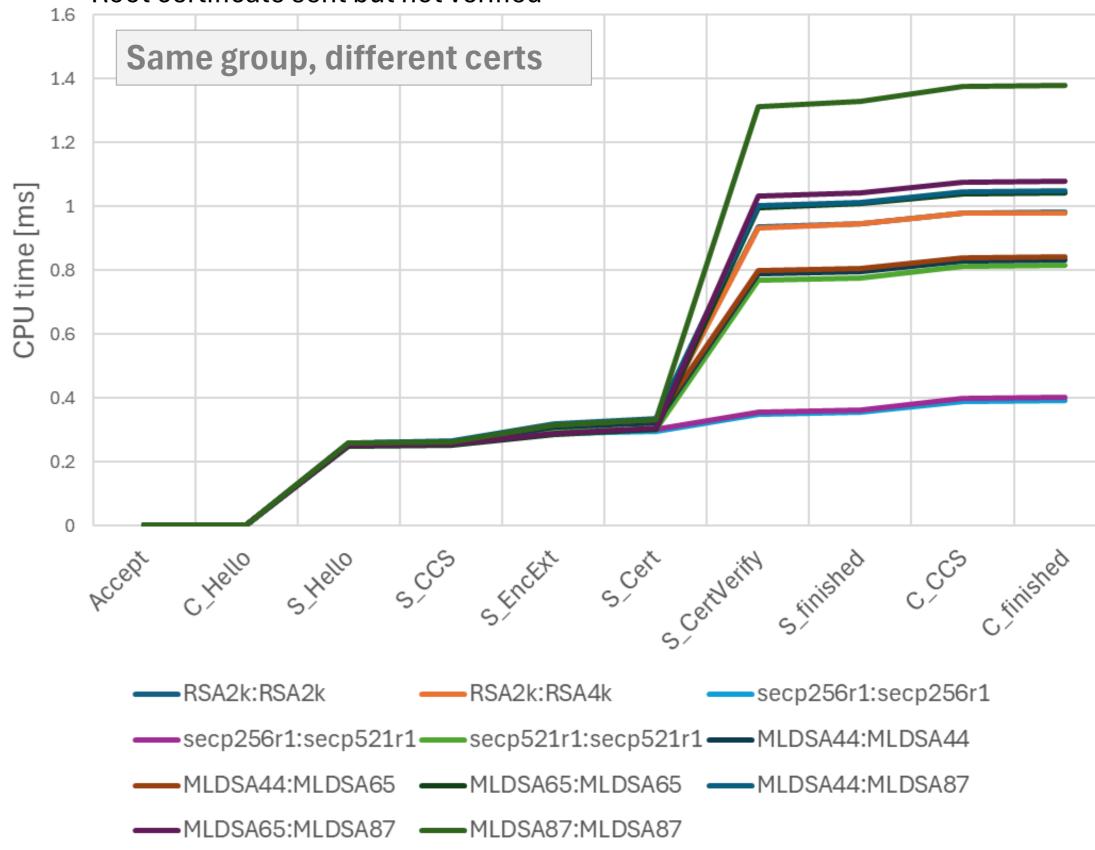
SecP384r1MLKEM1024

---MLKEM1024

**--**secp384r1

**X25519MLKEM768** 





Certificate notation: <Leaf cert>:<Root cert>

——SecP256r1MLKEM768 ——x448

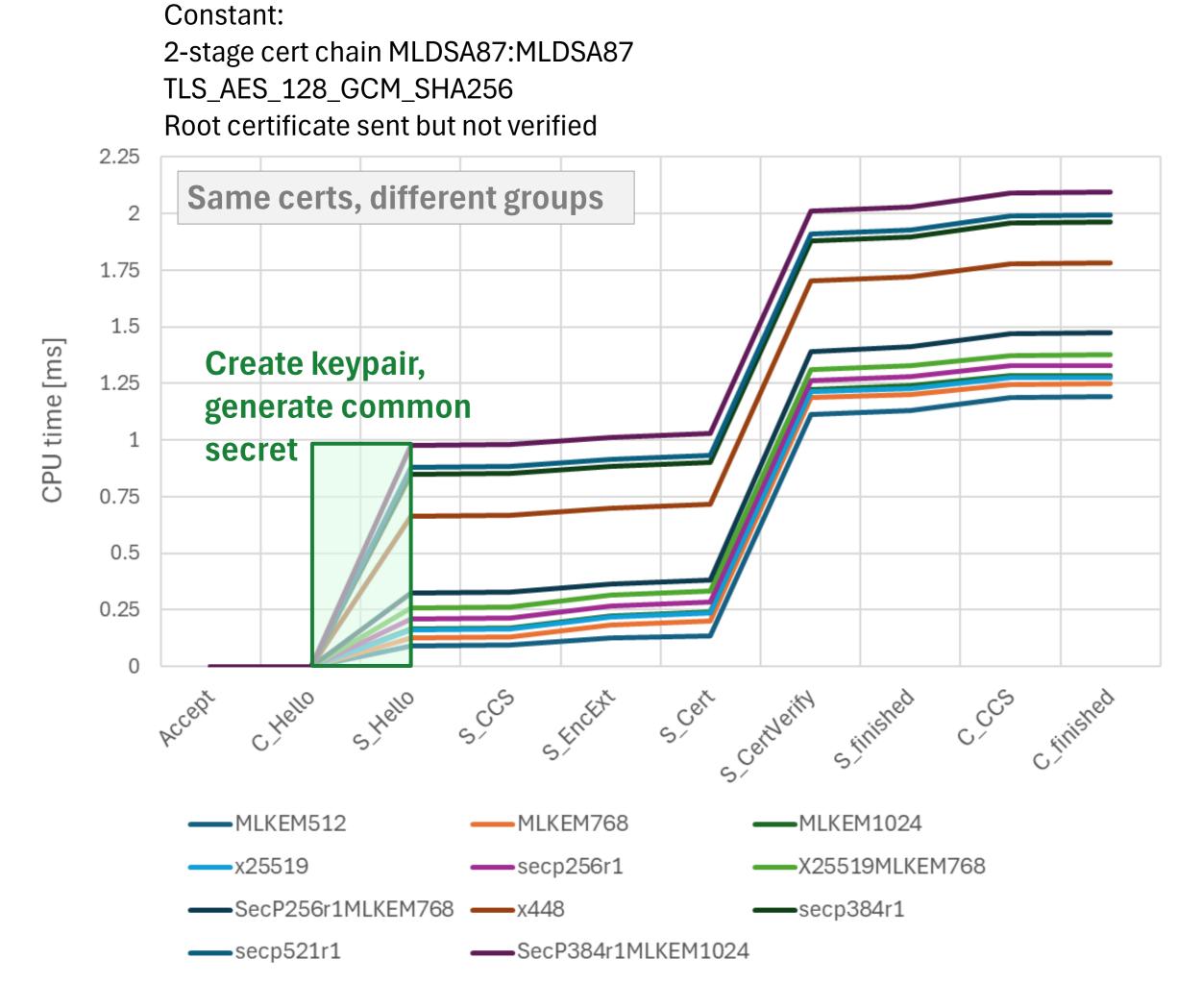
0.5

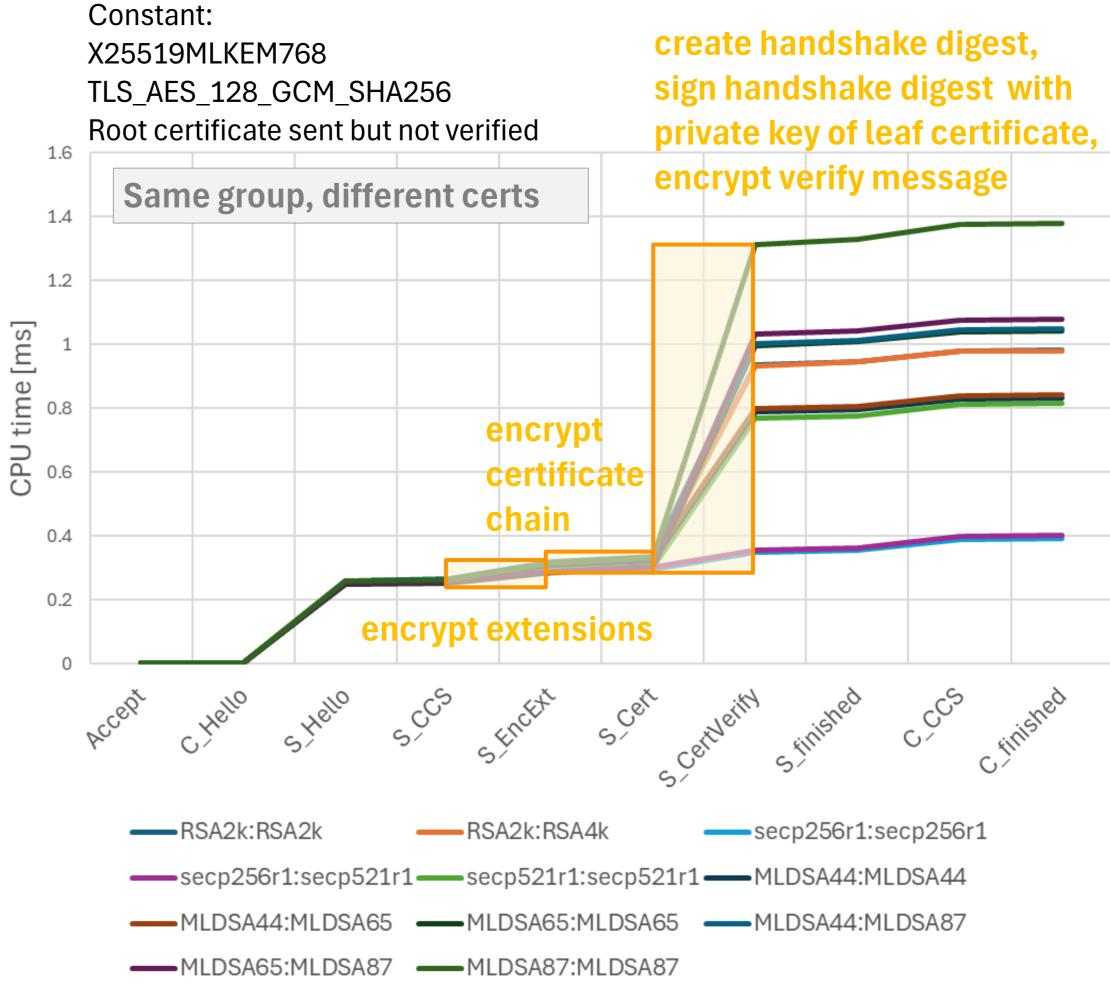
0.25

-MLKEM512

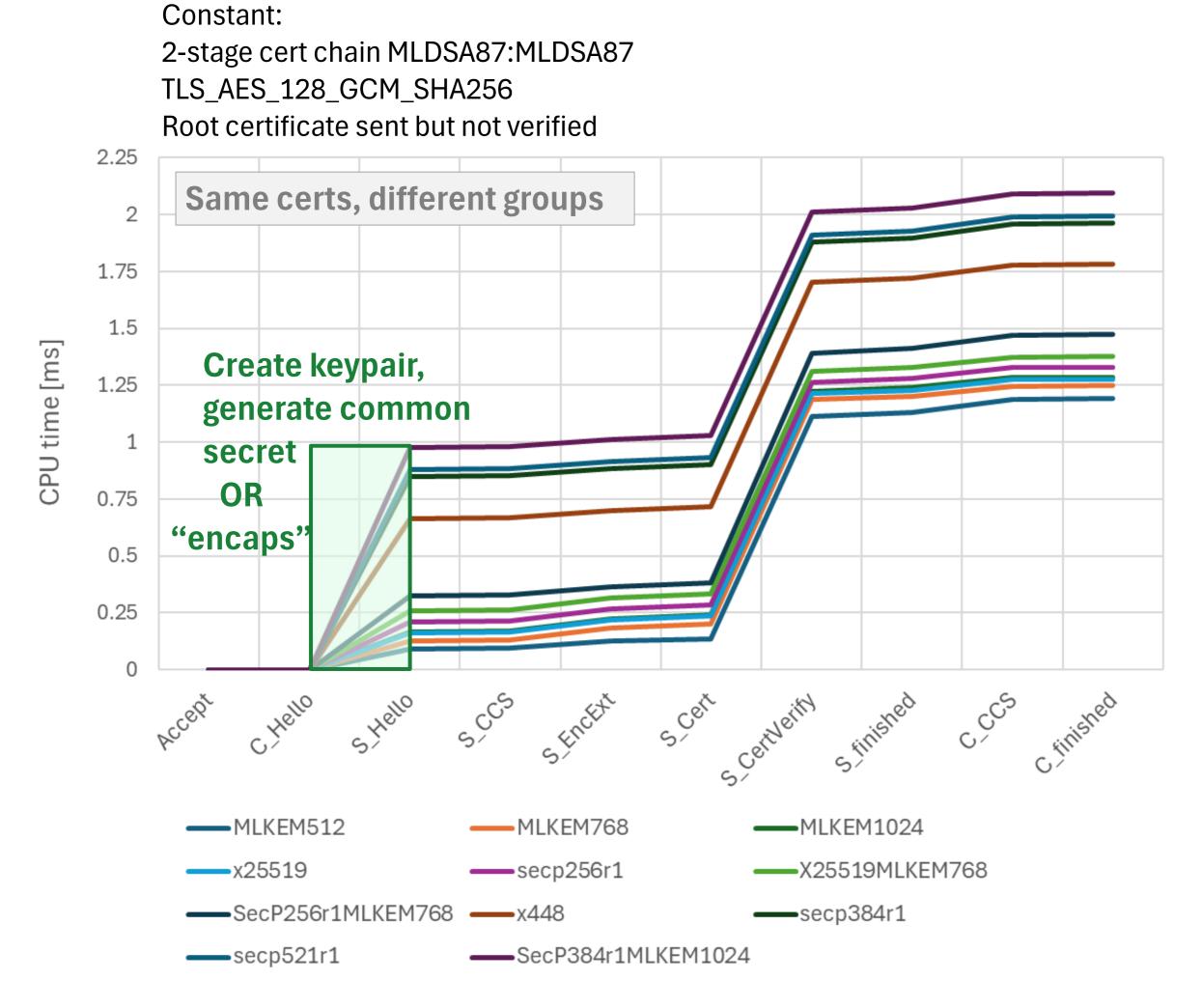
**--**x25519

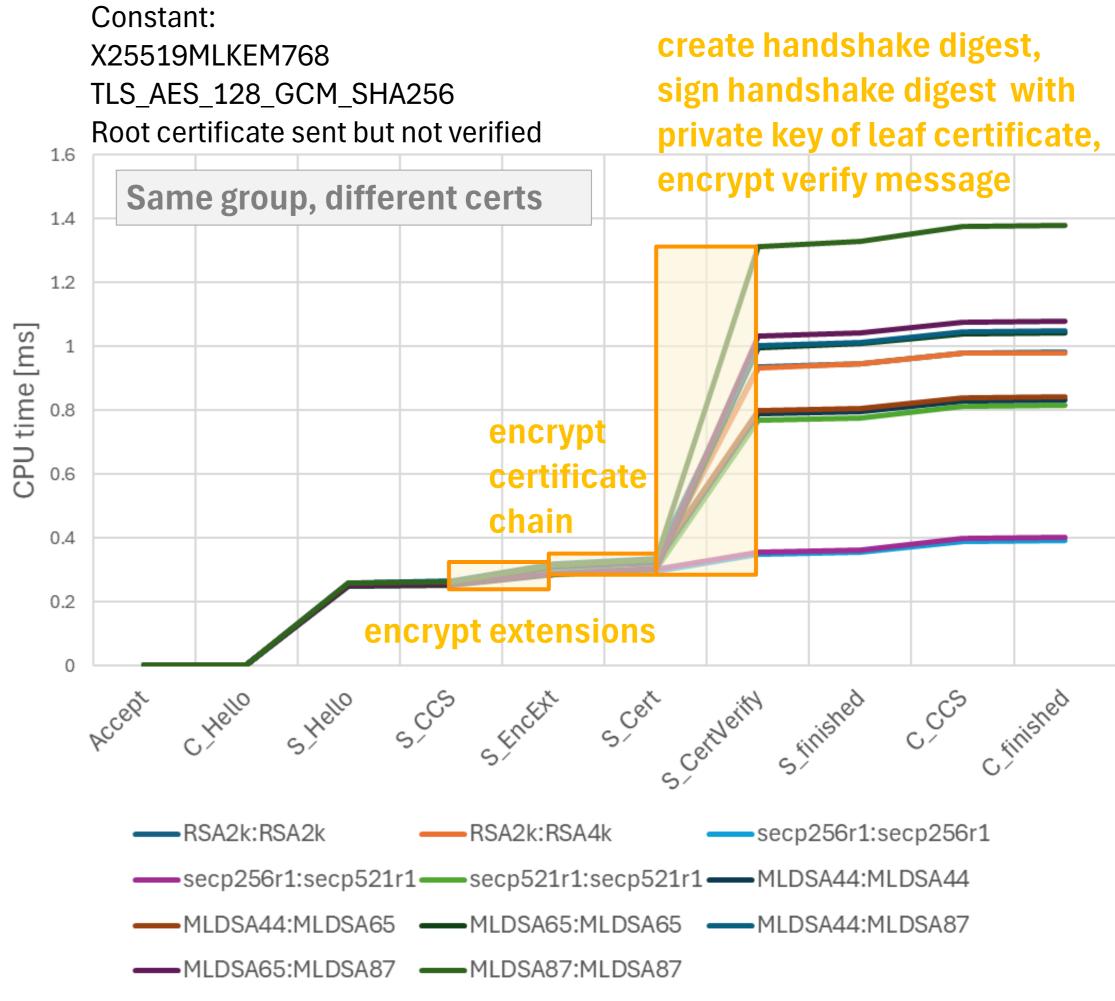
## Data extraction: SERVER





## Data extraction: SERVER





Certificate notation: <Leaf cert>:<Root cert>

## Measurement Setup

#### **Hardware**:

Bare metal rack server with Dual Intel(R) Xeon(R) "Ice Lake" Gold 6338 CPU @ 2.00GHz, 512 GB  $\rightarrow$  Total 64C/128T C-states & "Turbo" OFF

#### Setup:

TLS client & TLS server in c-code

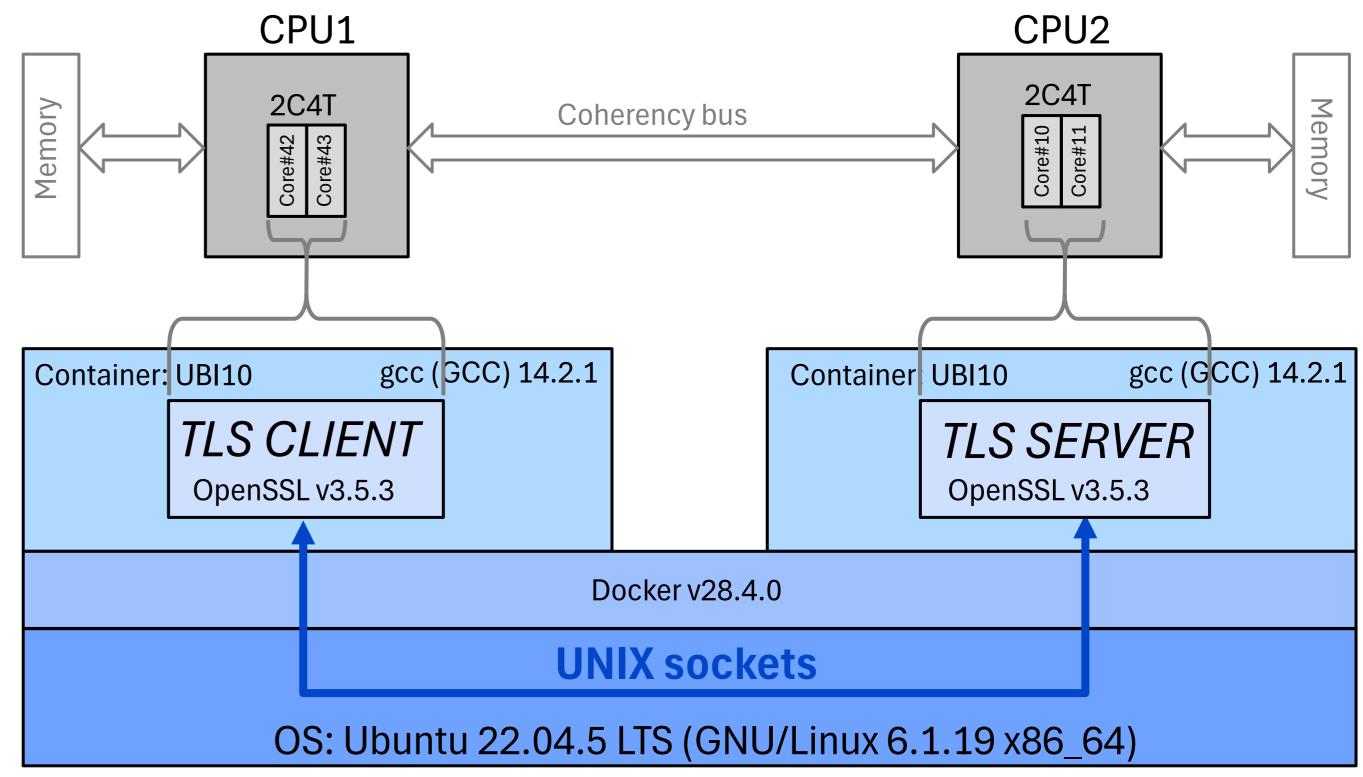
- on same bare metal computer
- each running on 2C4T (CPU pinning to two neighboring cores)
- each on its 'own' CPU chip
- each in its own Docker container
- sequential single TLS sessions
- communication via UNIX sockets
- single supported group and single keyshare in ClientHello → no HRR, no encExt

#### **OpenSSL** build options:

--prefix=/usr --with-rand-seed=rdcpu,os -DFAST\_PCLMUL -DO3 -DECP\_NISTZ256\_ASM -DX25519\_ASM -DCHACHA\_ASM -POLY1305\_ASM -DOPENSSL\_BN\_ASM\_MONT -DSHA1\_ASM -DSHA256\_ASM -DSHA512\_ASM -DKECCAK1600\_ASM -DAES\_ASM -DVPAES\_ASM enable-tls1\_3 enable-ec\_nistp\_64\_gcc\_128 no-ssl no-tls1 no-tls1\_1 no-afalgeng no-tests shared threads -lm -mrdrnd

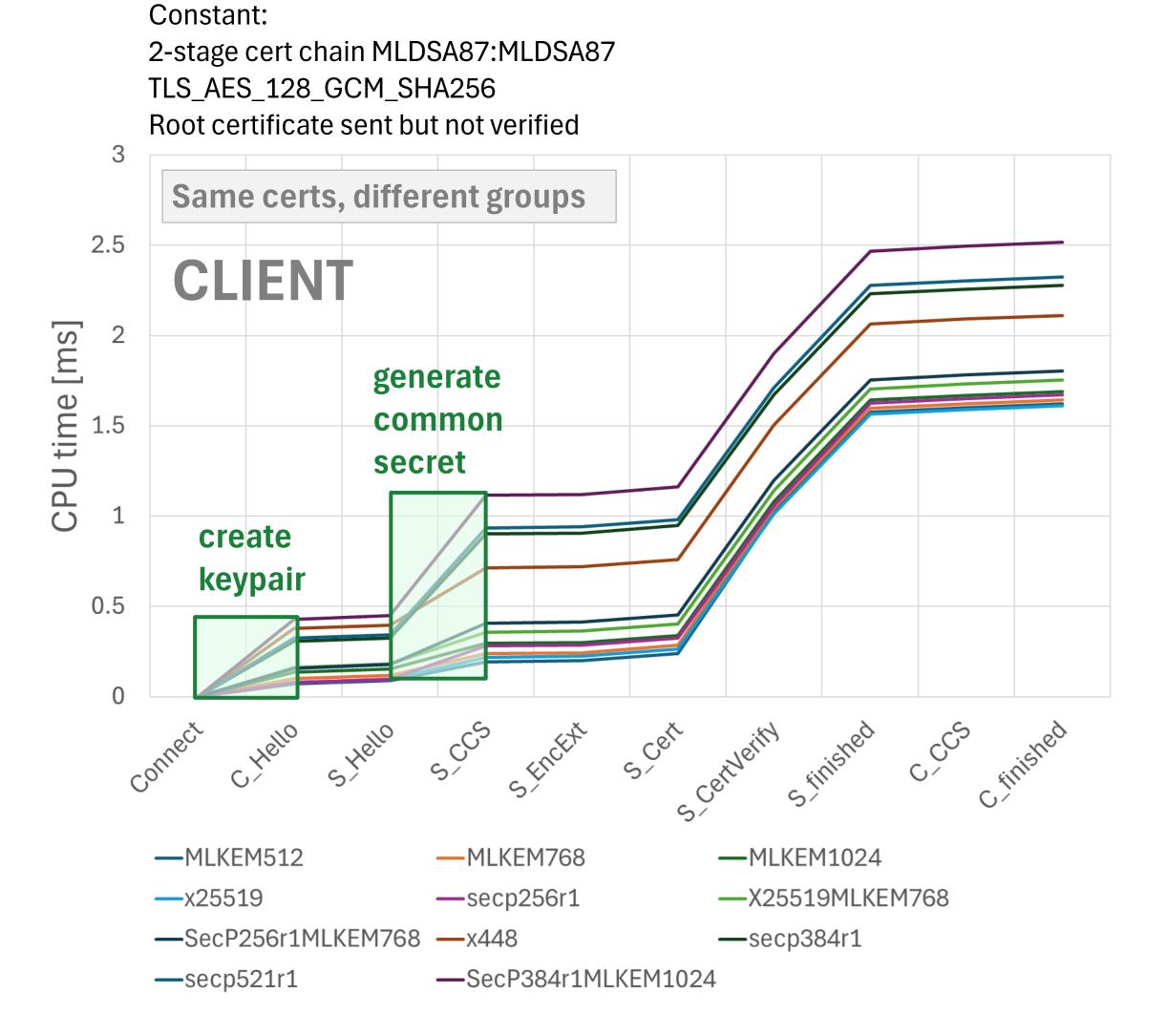
#### **Experiment duration**:

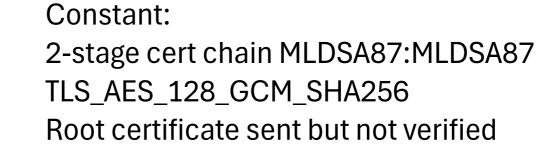
Each experiment with 10'000 warmup loops and 100'000 measurement loops Measurement results show the 50% quantile (median) unless otherwise noted Repeatability verified to be within  $\pm 1\%$ 

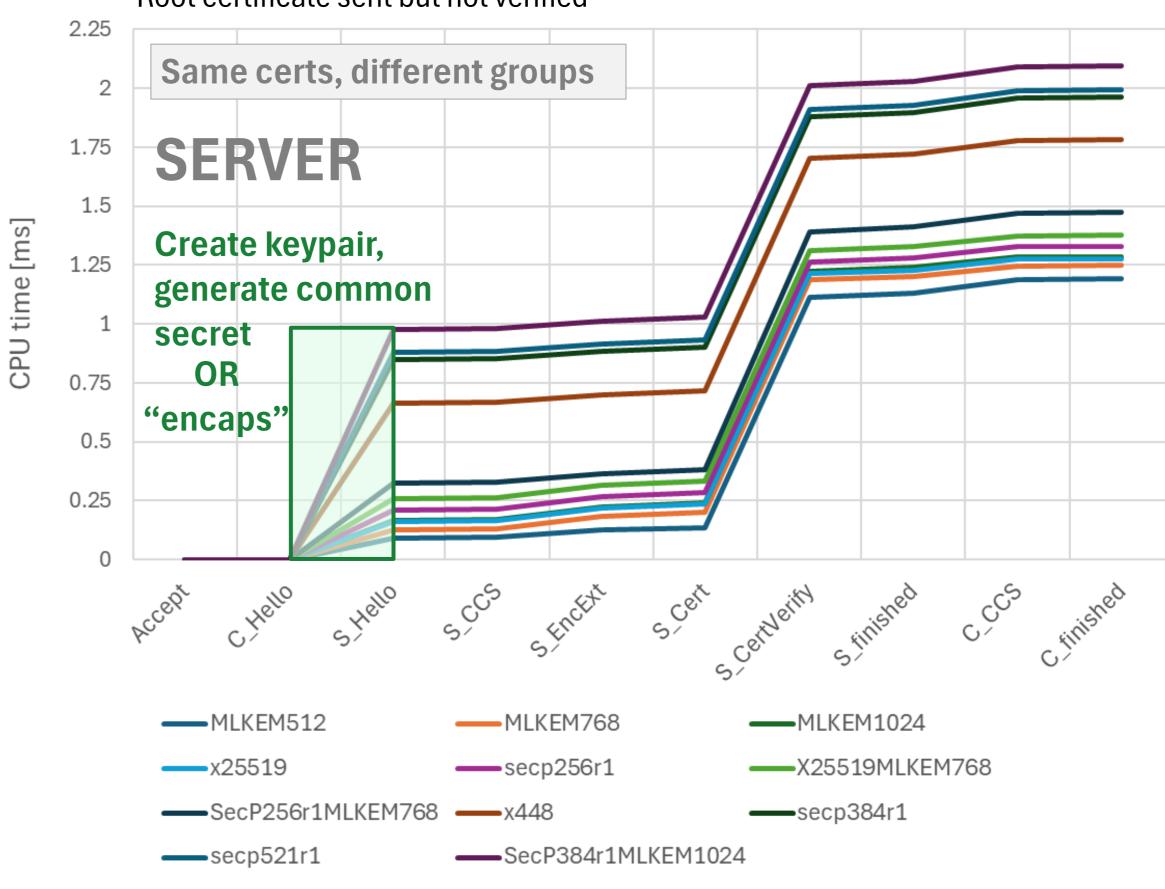


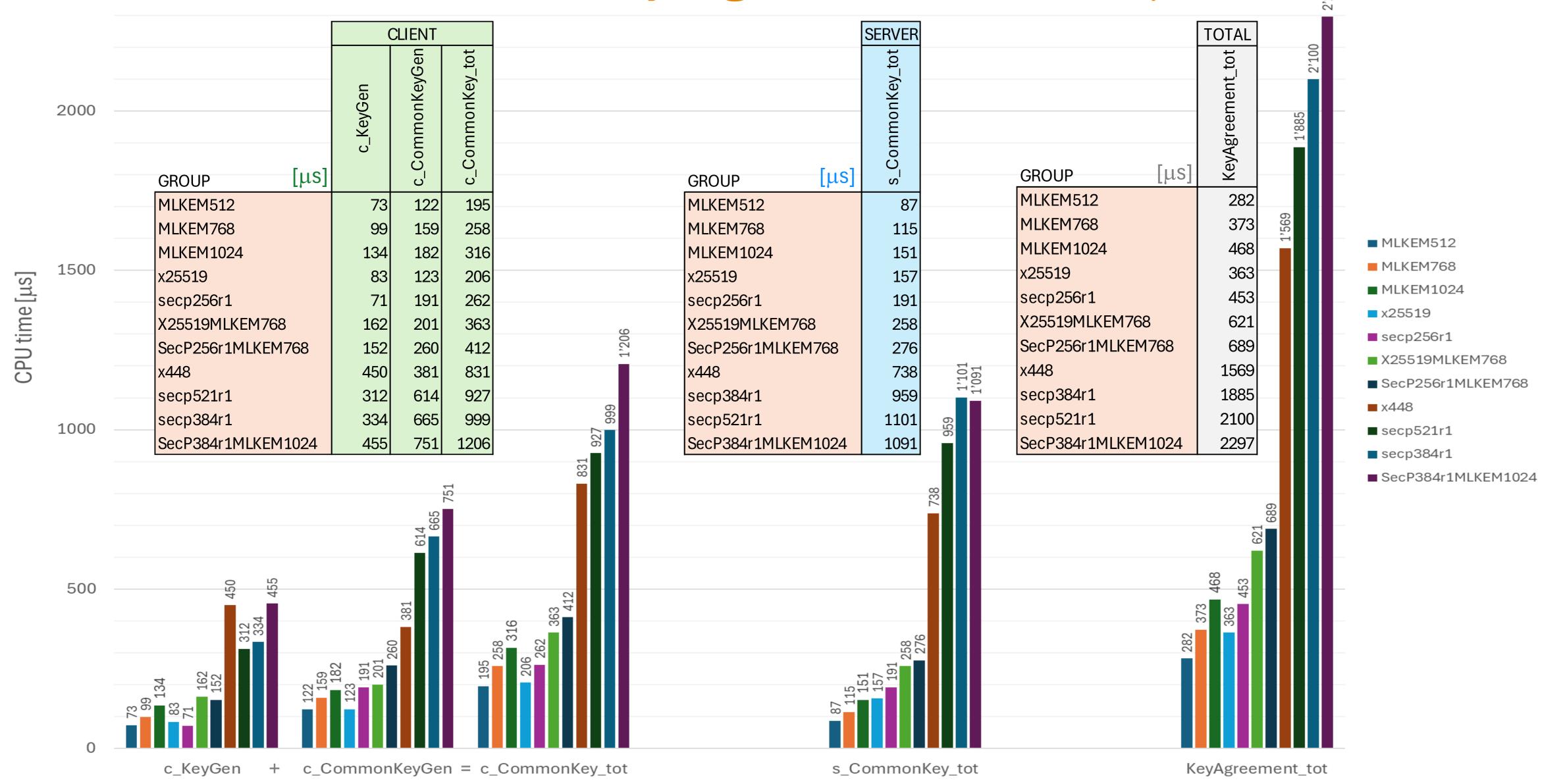
Remark: We did not see any impact on CPU times from using UNIX sockets or TCP via localhost

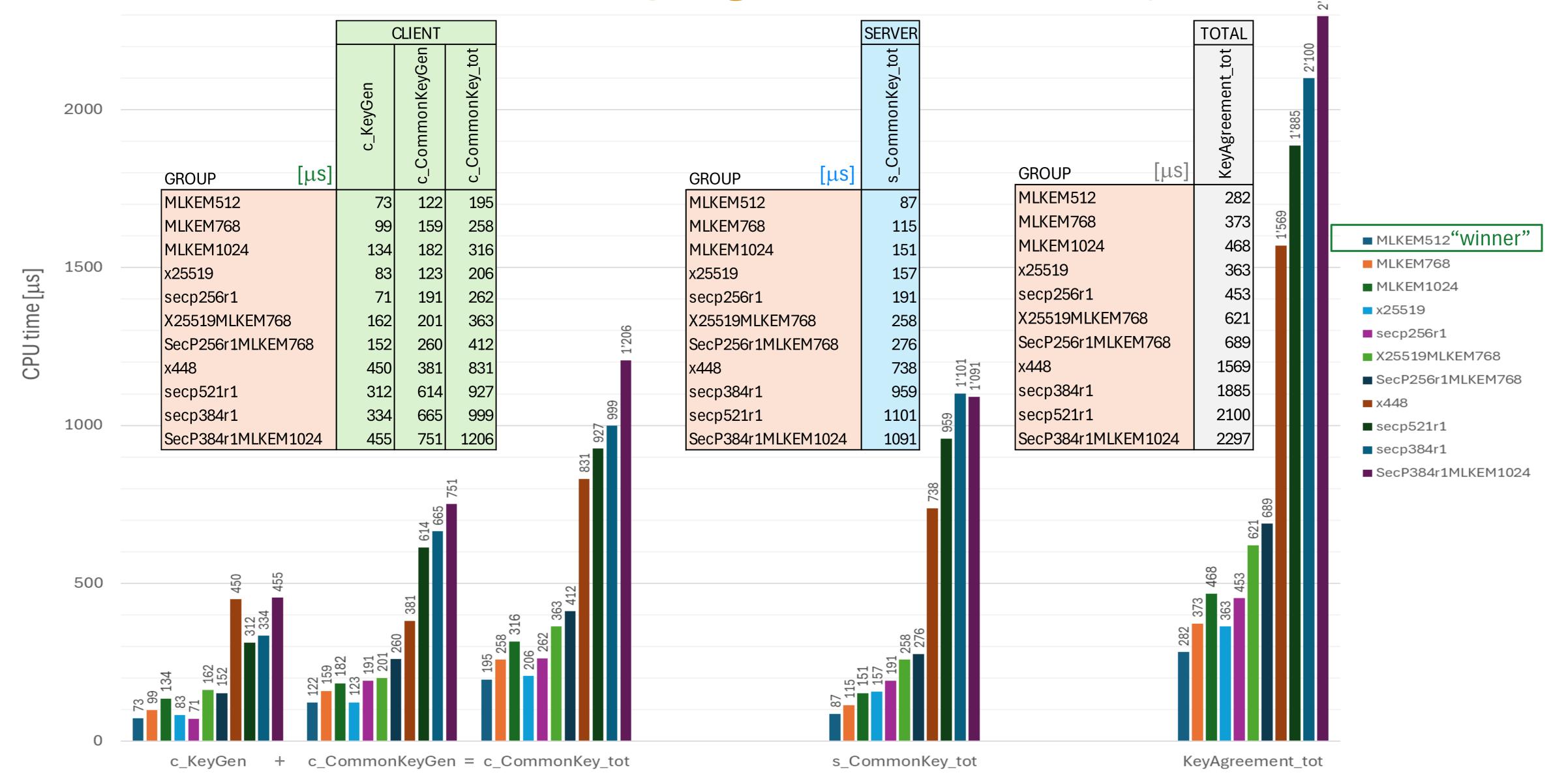
# Measured Results: Key Agreement

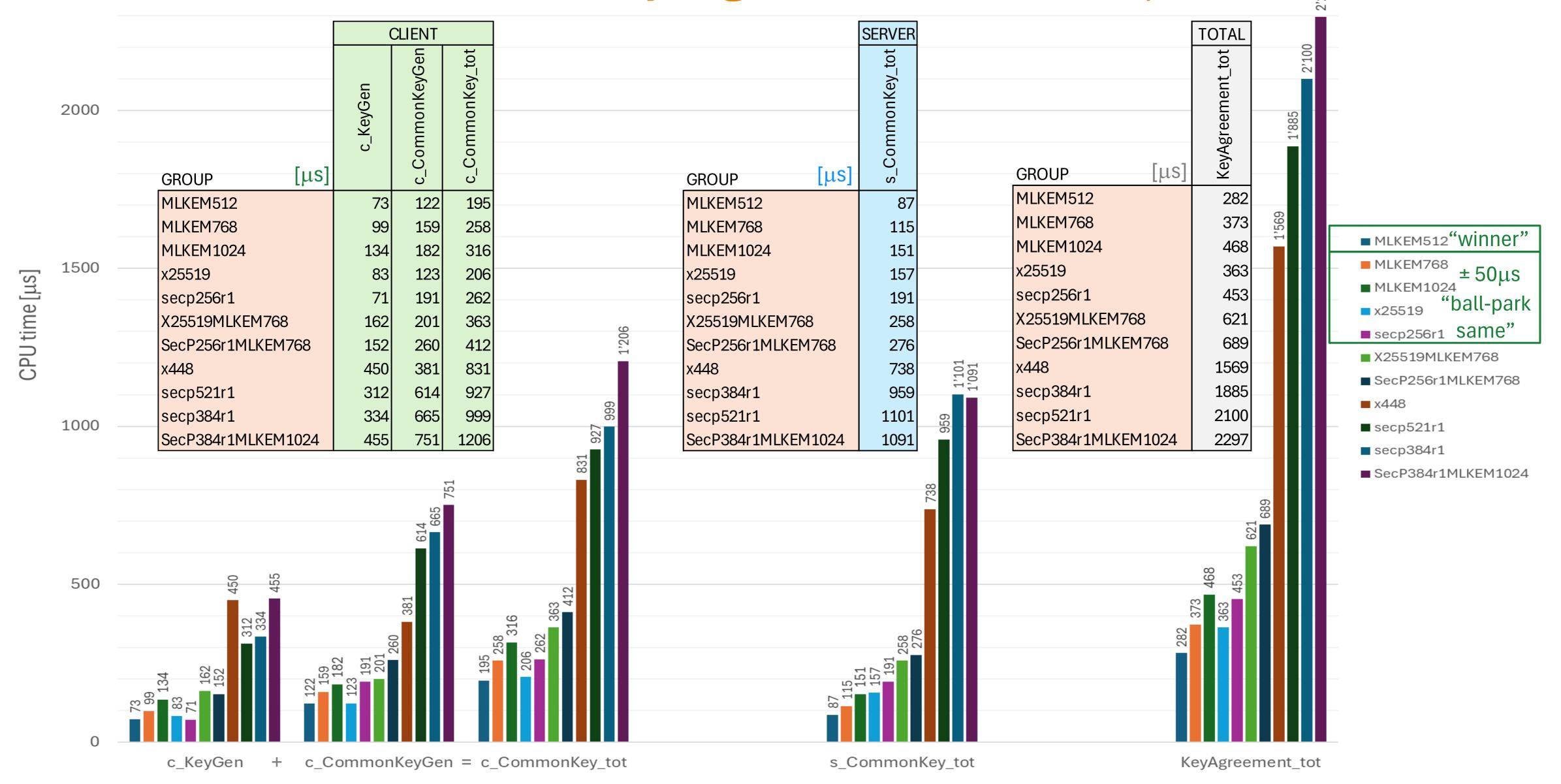


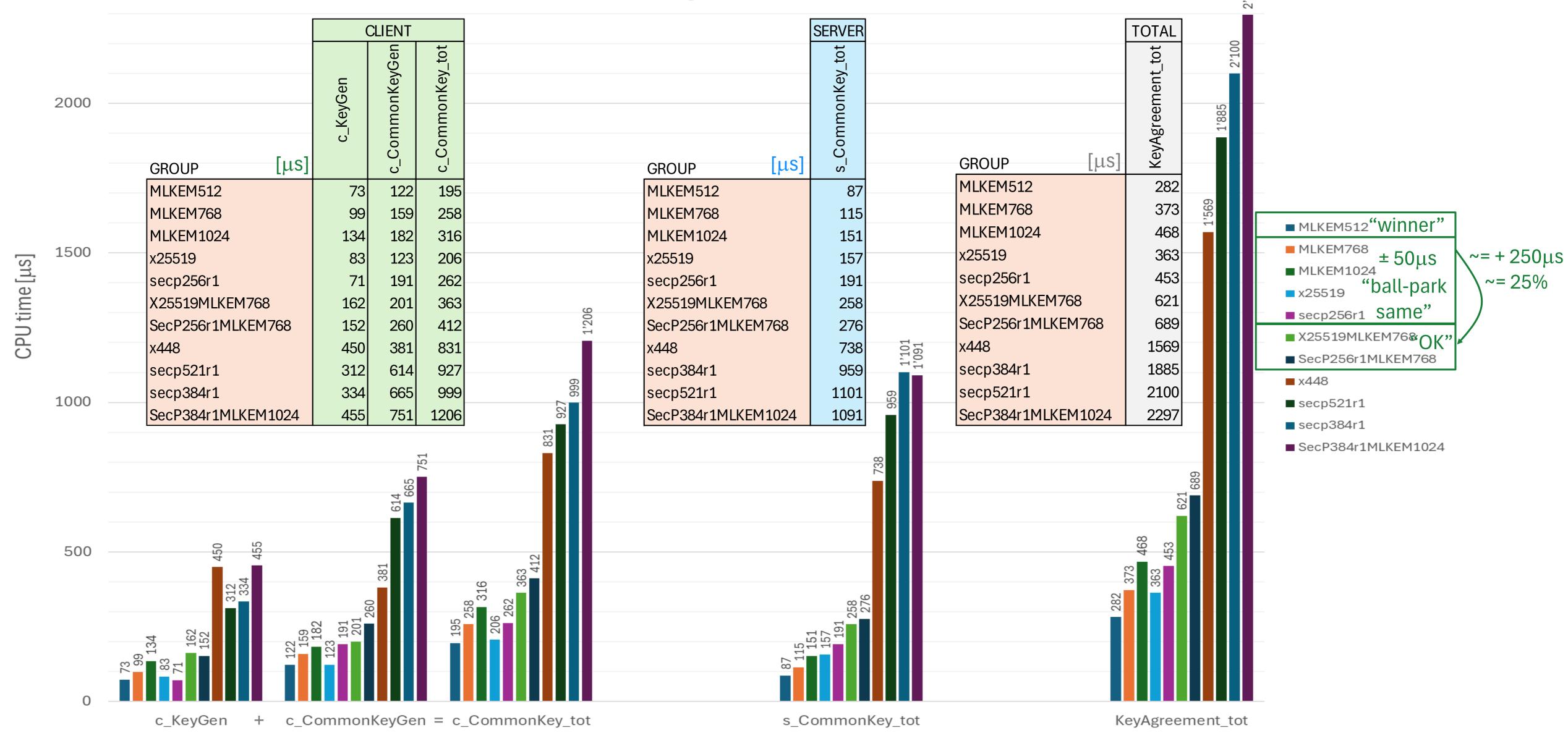


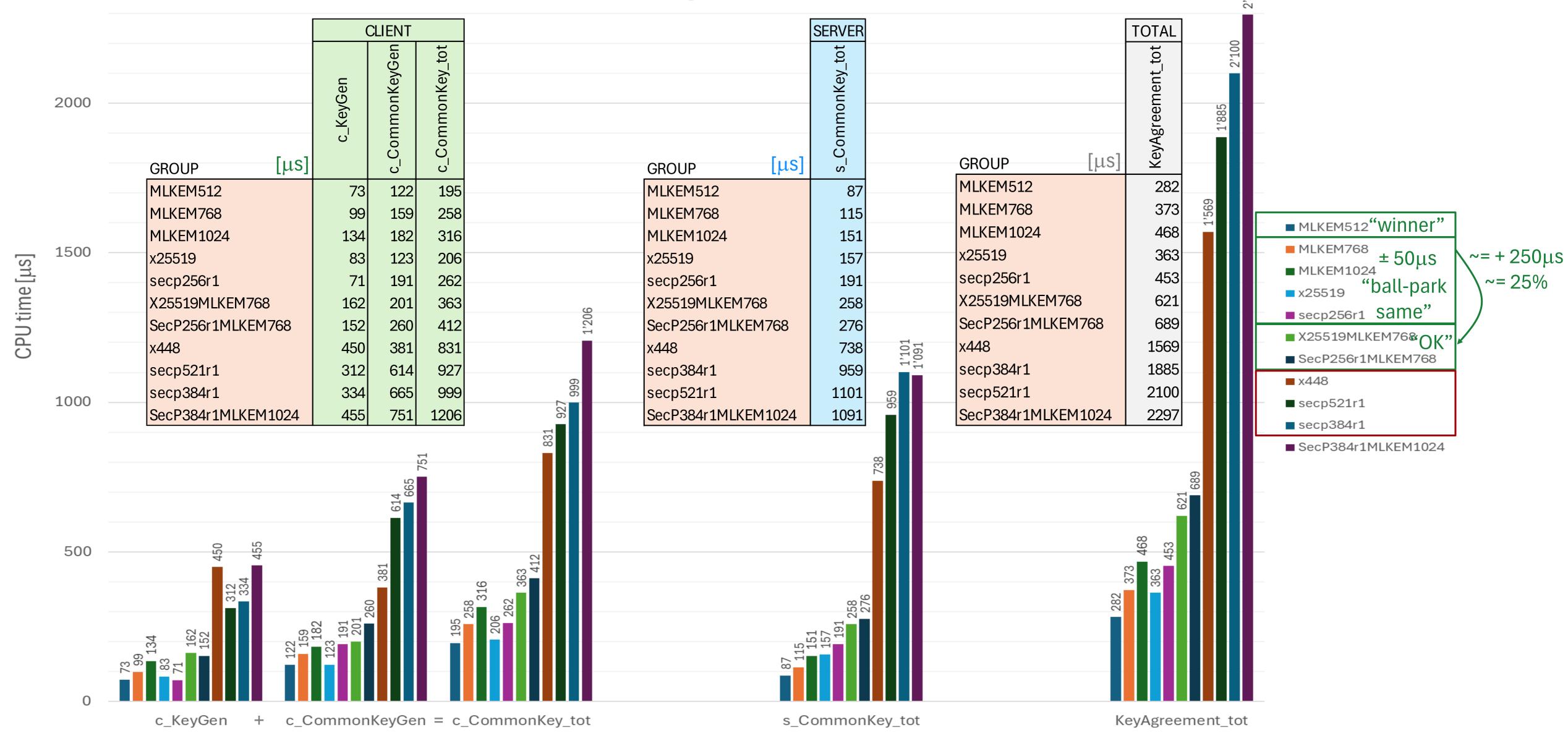


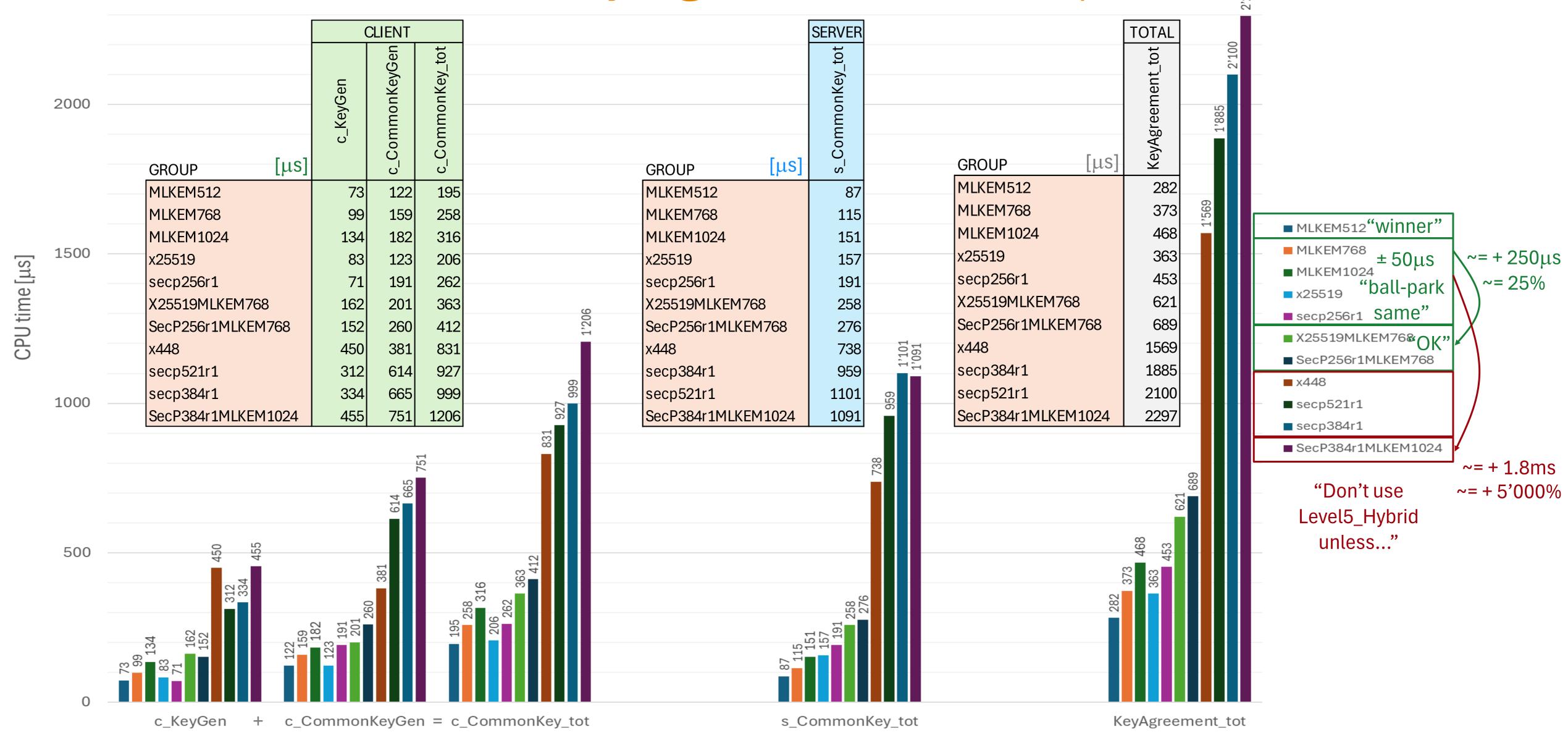


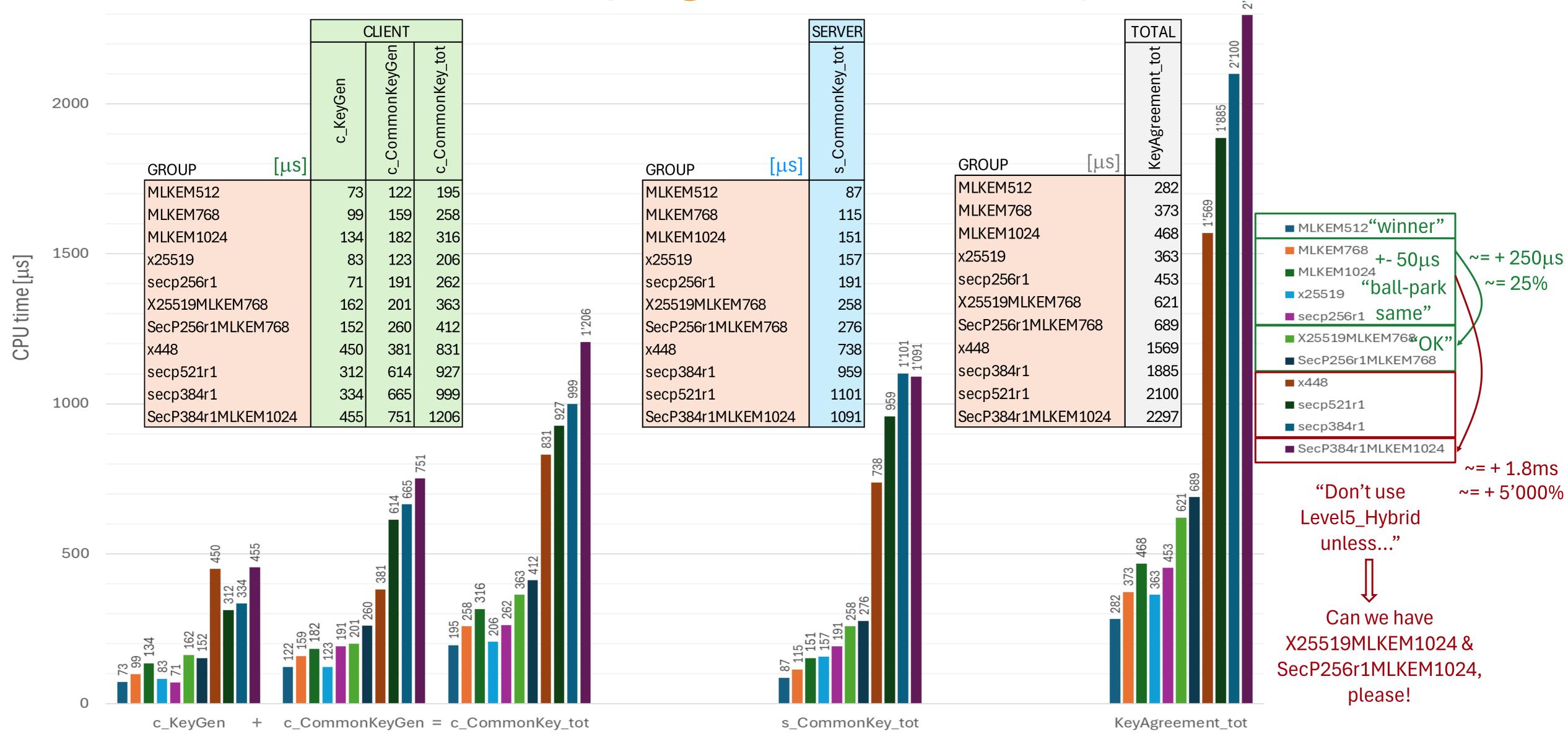


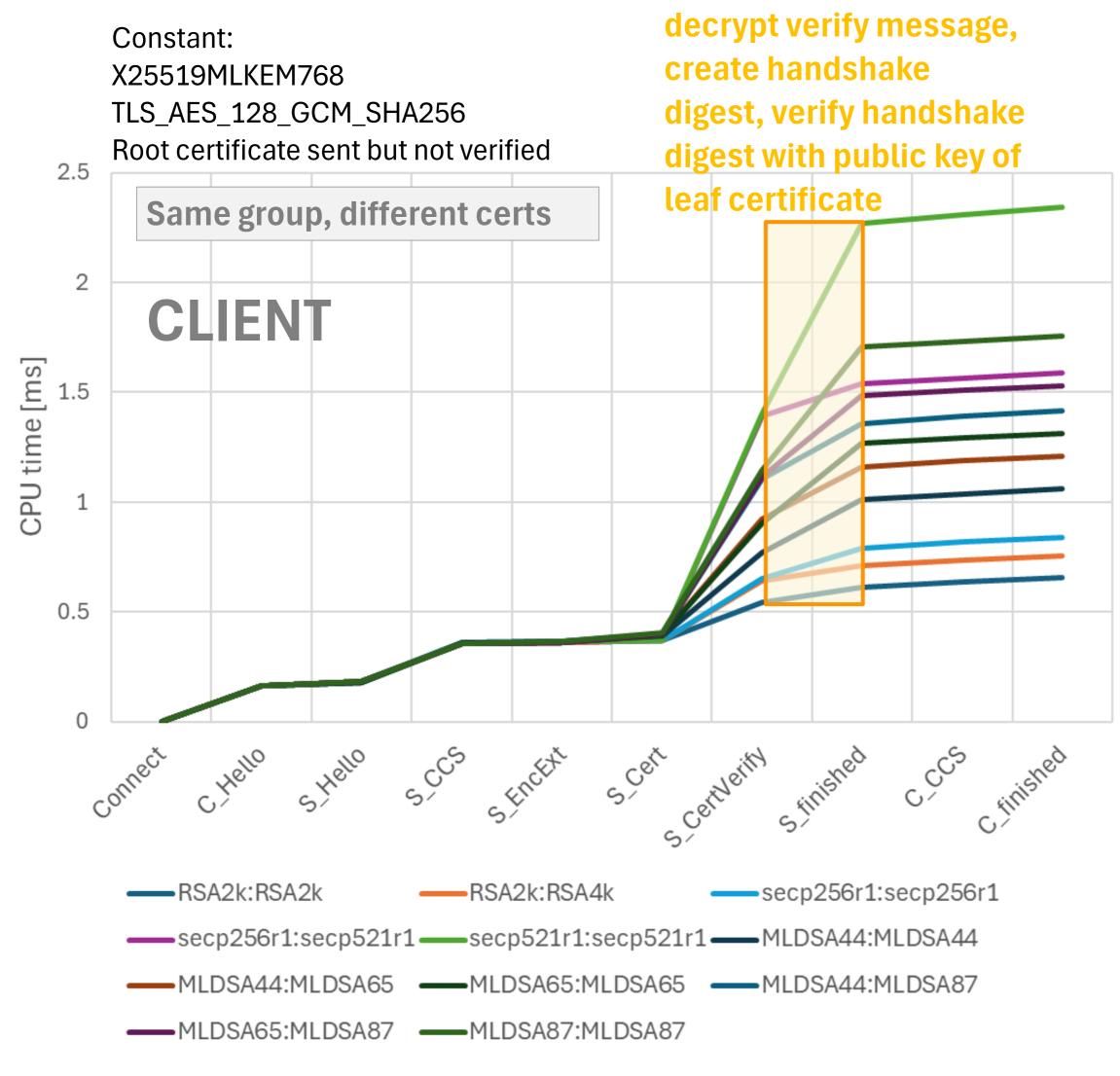




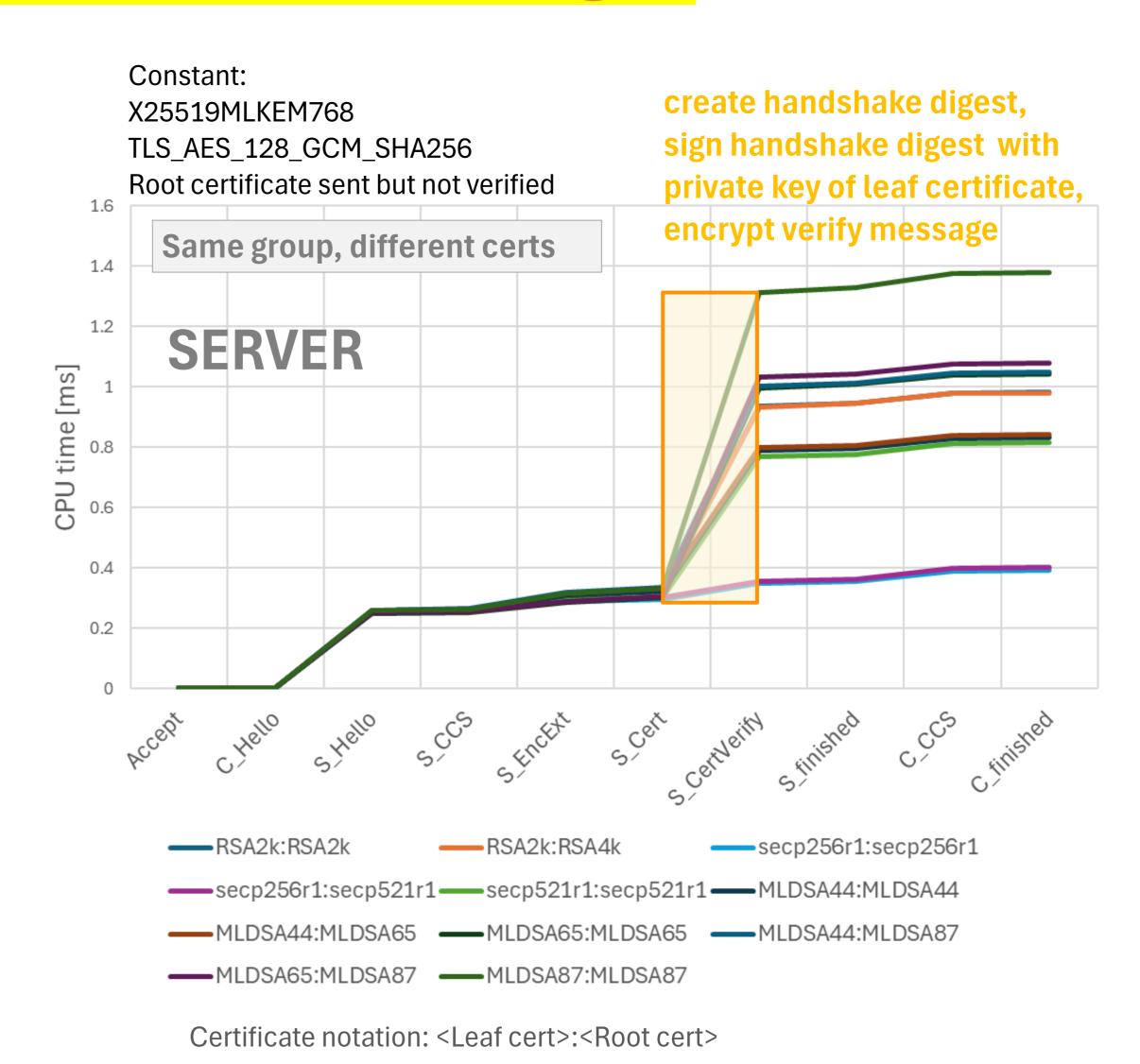




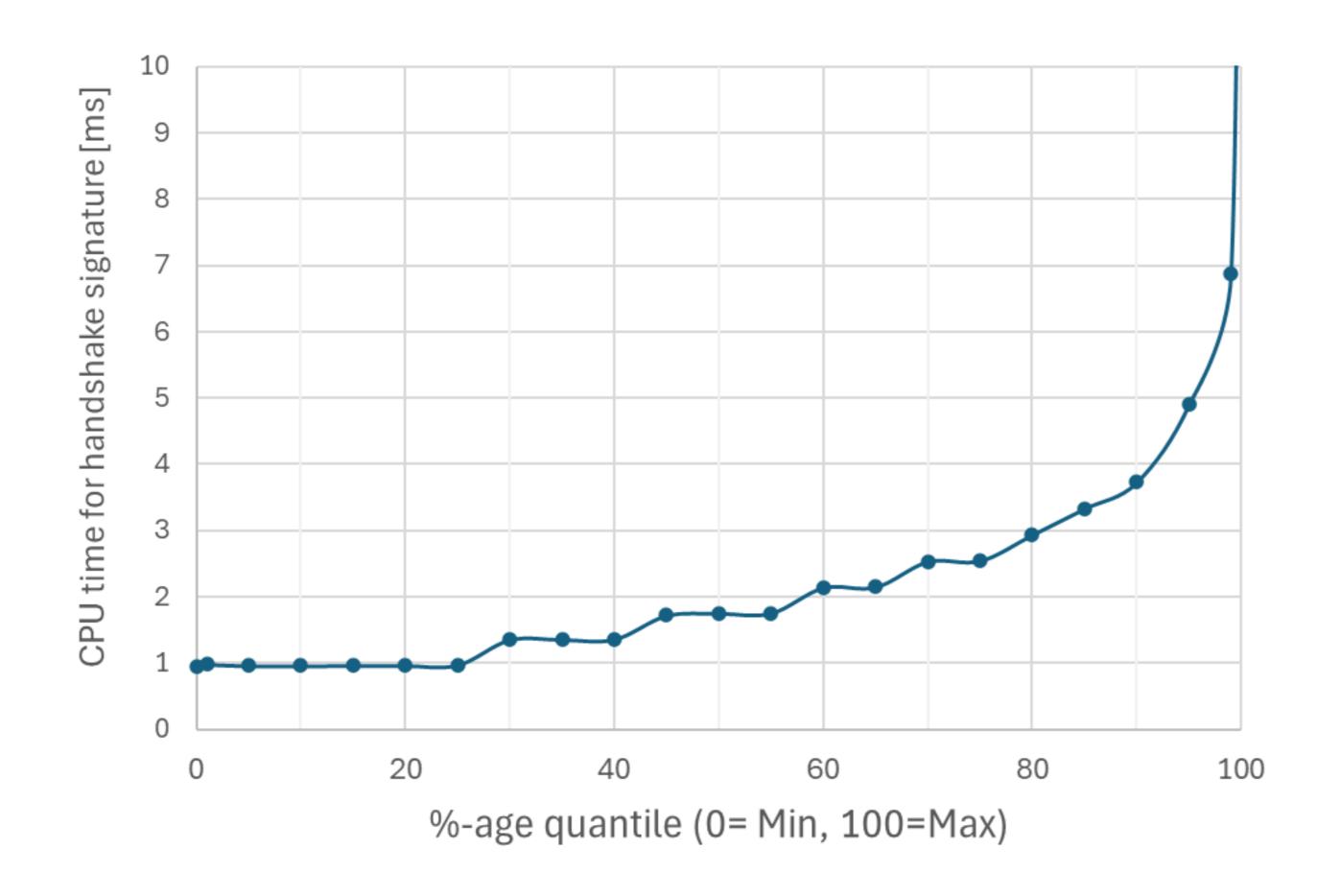




Certificate notation: <Leaf cert>:<Root cert>



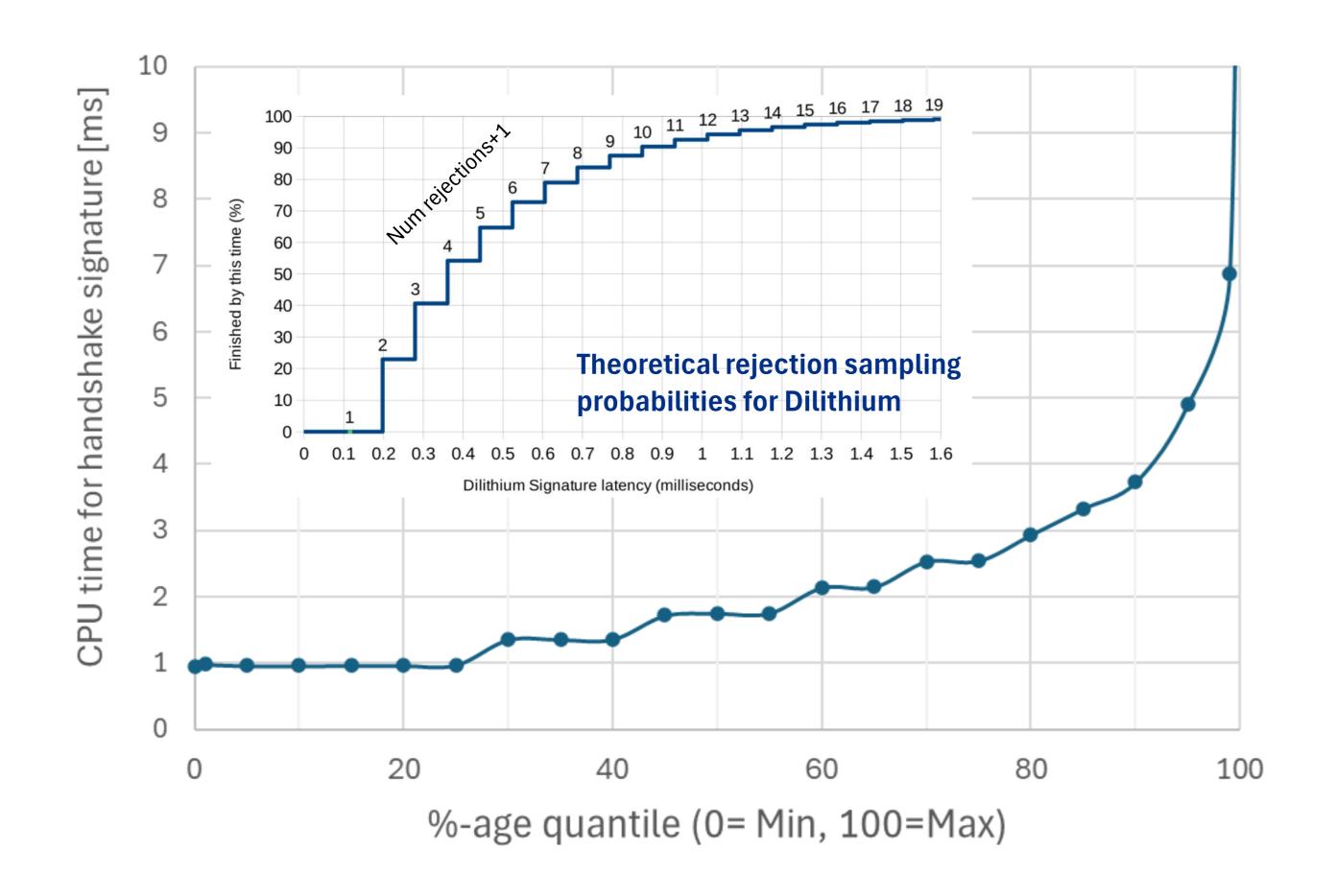
Group X25519 2-stage cert chain MLDSA87:MLDSA87 TLS\_AES\_128\_GCM\_SHA256 Root certificate sent but not verified



#### MLDSA rejection sampling leads to unbound latency variations and hence to tail latencies

Rejection sampling is a security mechanism that ensures signatures are statistically close to secret-independent and corrects for potential leakage of secret information during signature generation.

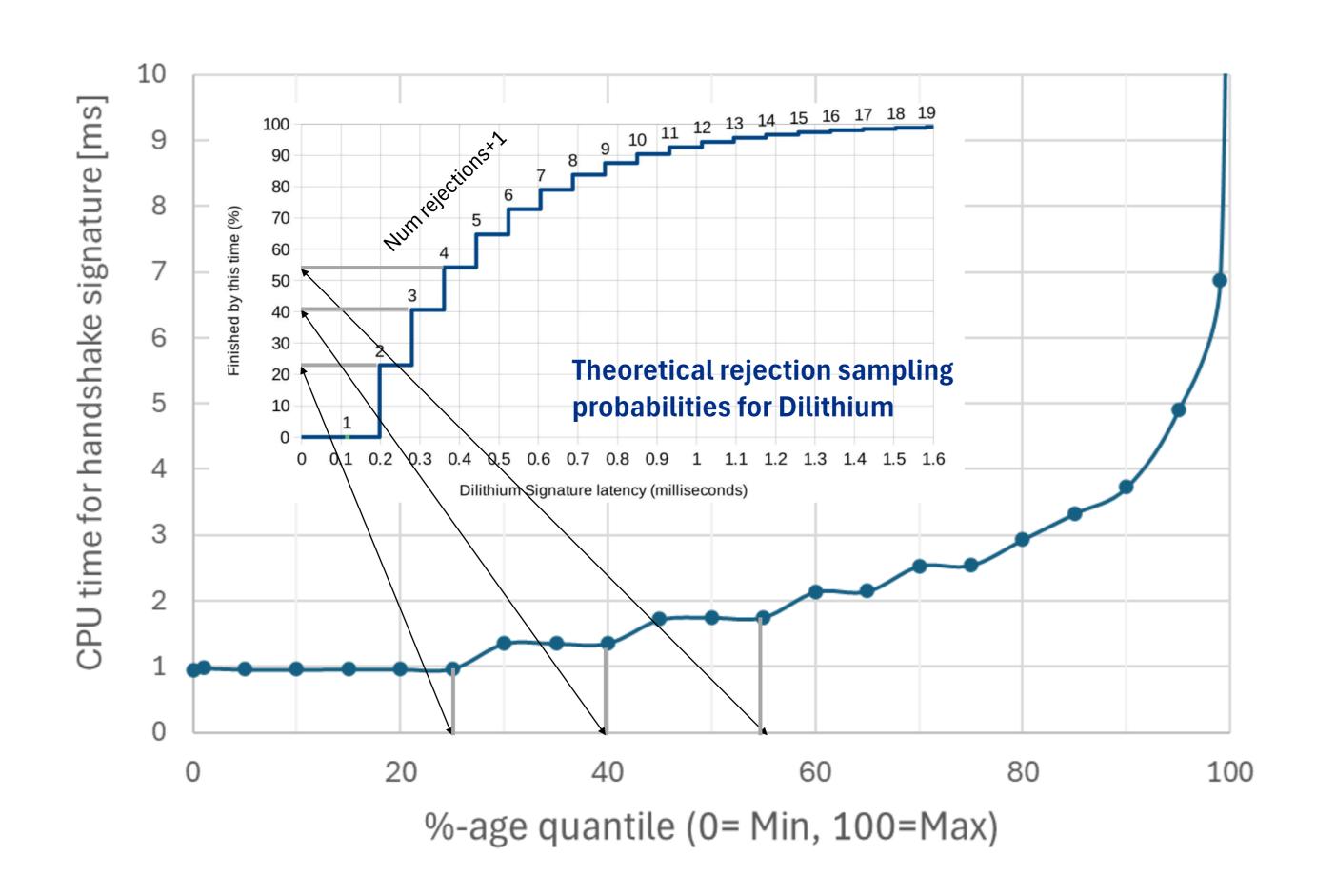
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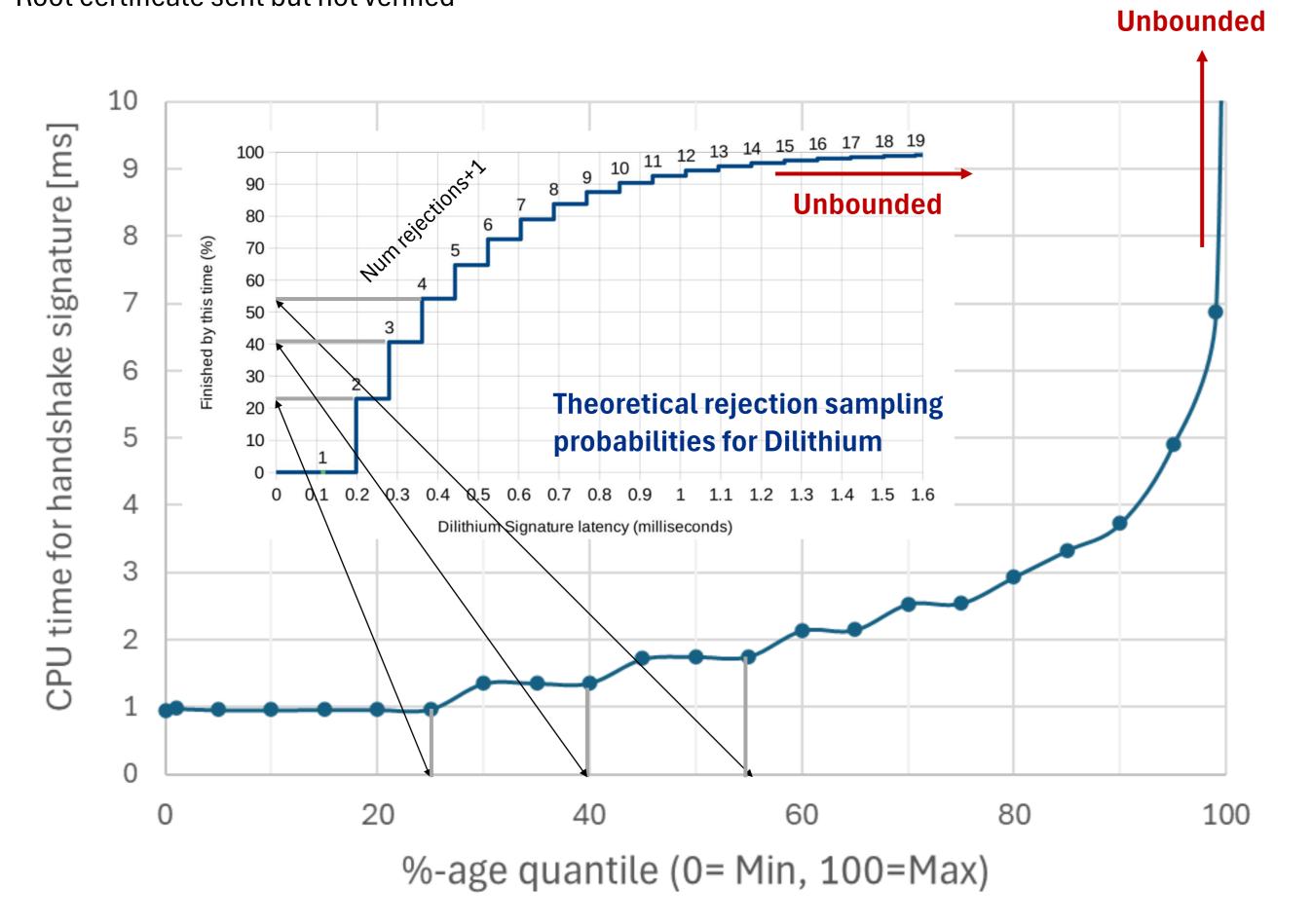
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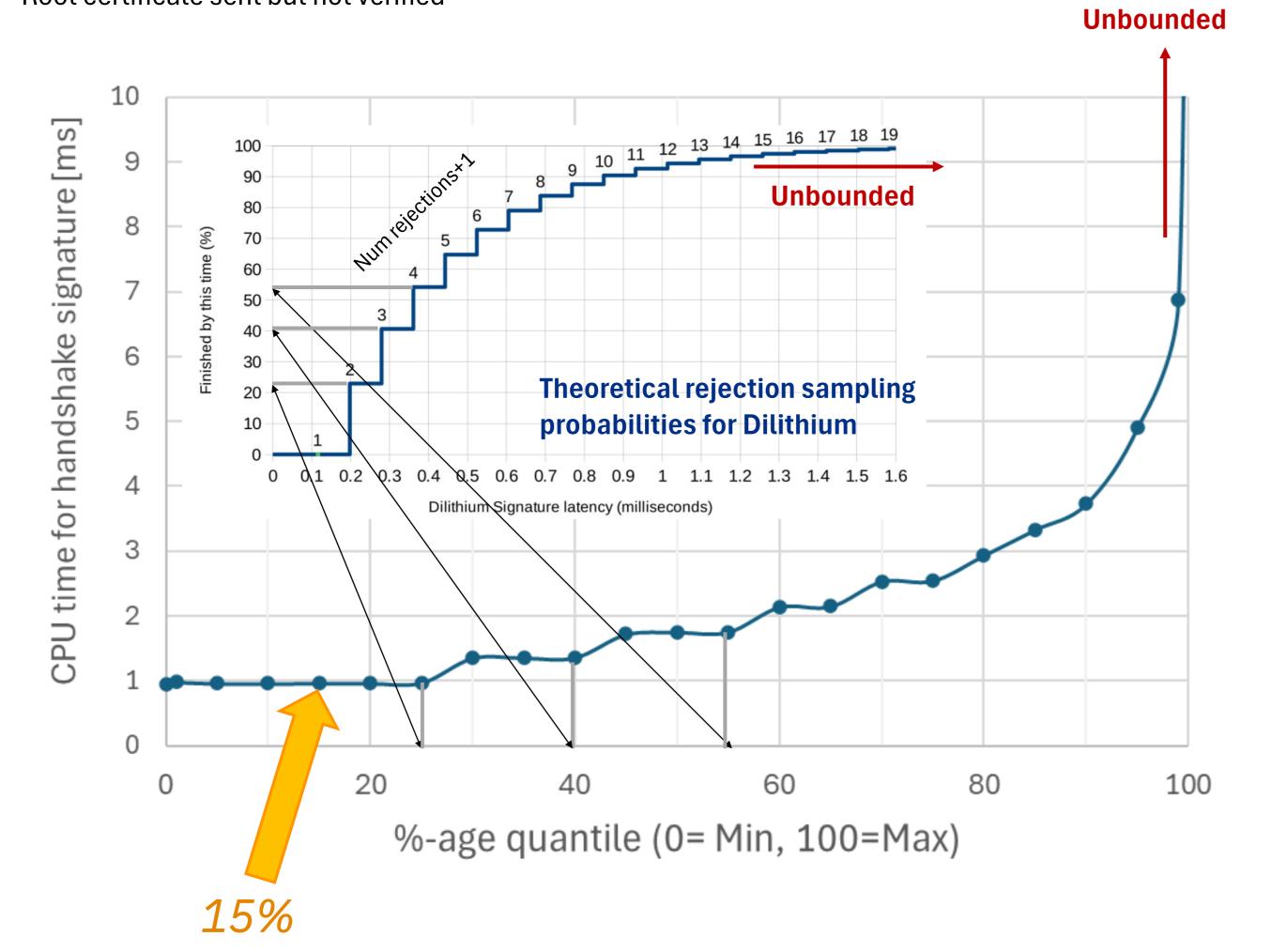
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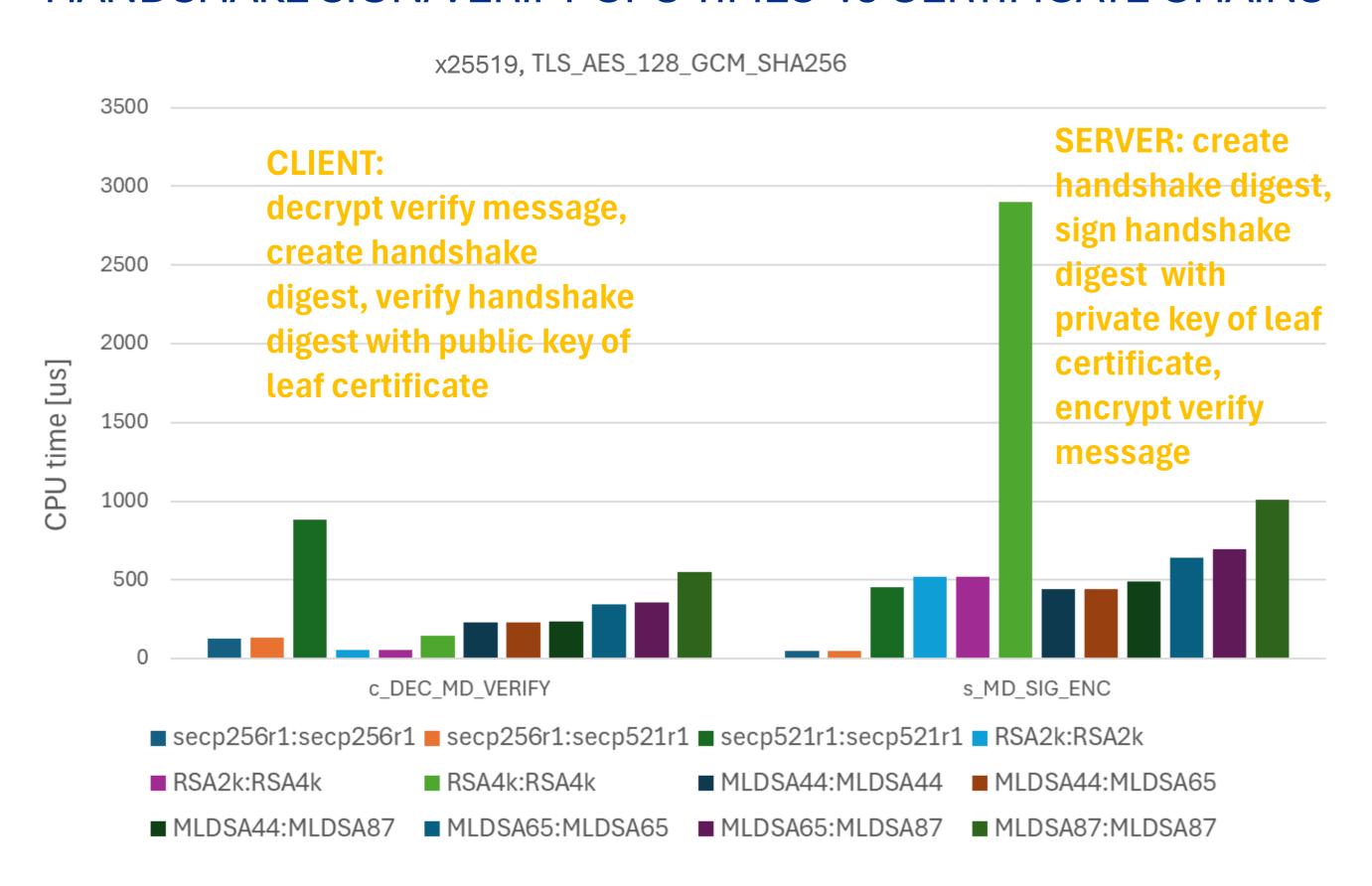


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Rejection sampling is a security mechanism that ensures signatures are statistically close to secret-independent and corrects for potential leakage of secret information during signature generation.

- Can we have sign/verify algorithms without rejection sampling, please!
- → We show results for the 15% quantile

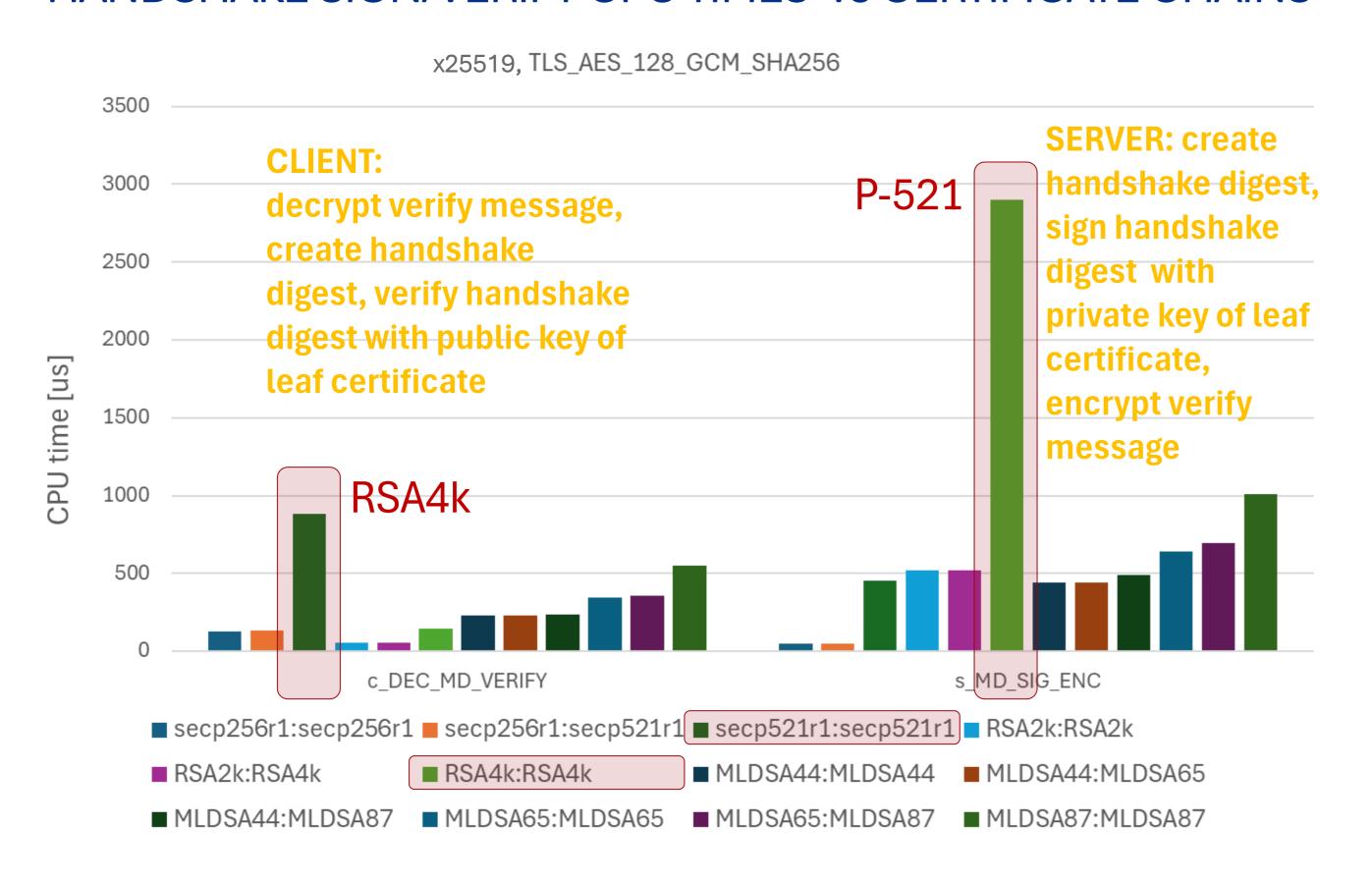
#### HANDSHAKE SIGN/VERIFY CPU TIMES vs CERTIFICATE CHAINS



Only the crypto algorithm from the leaf certificate is used to sign/verify

Ciphers have a small influence ( $\pm$  50  $\mu$ s), hence we show results only for TLS\_AES\_128\_GCM\_SHA256

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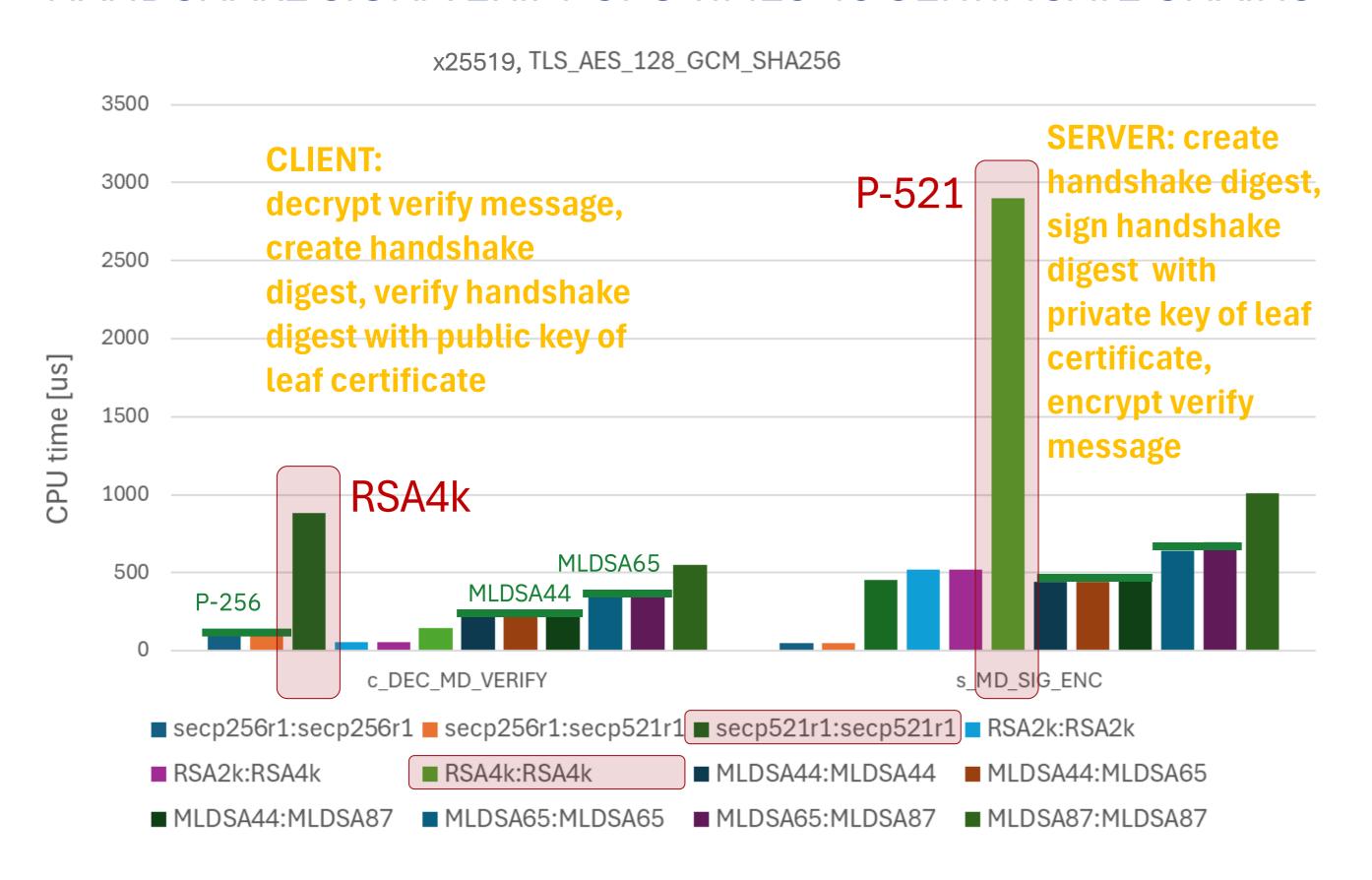


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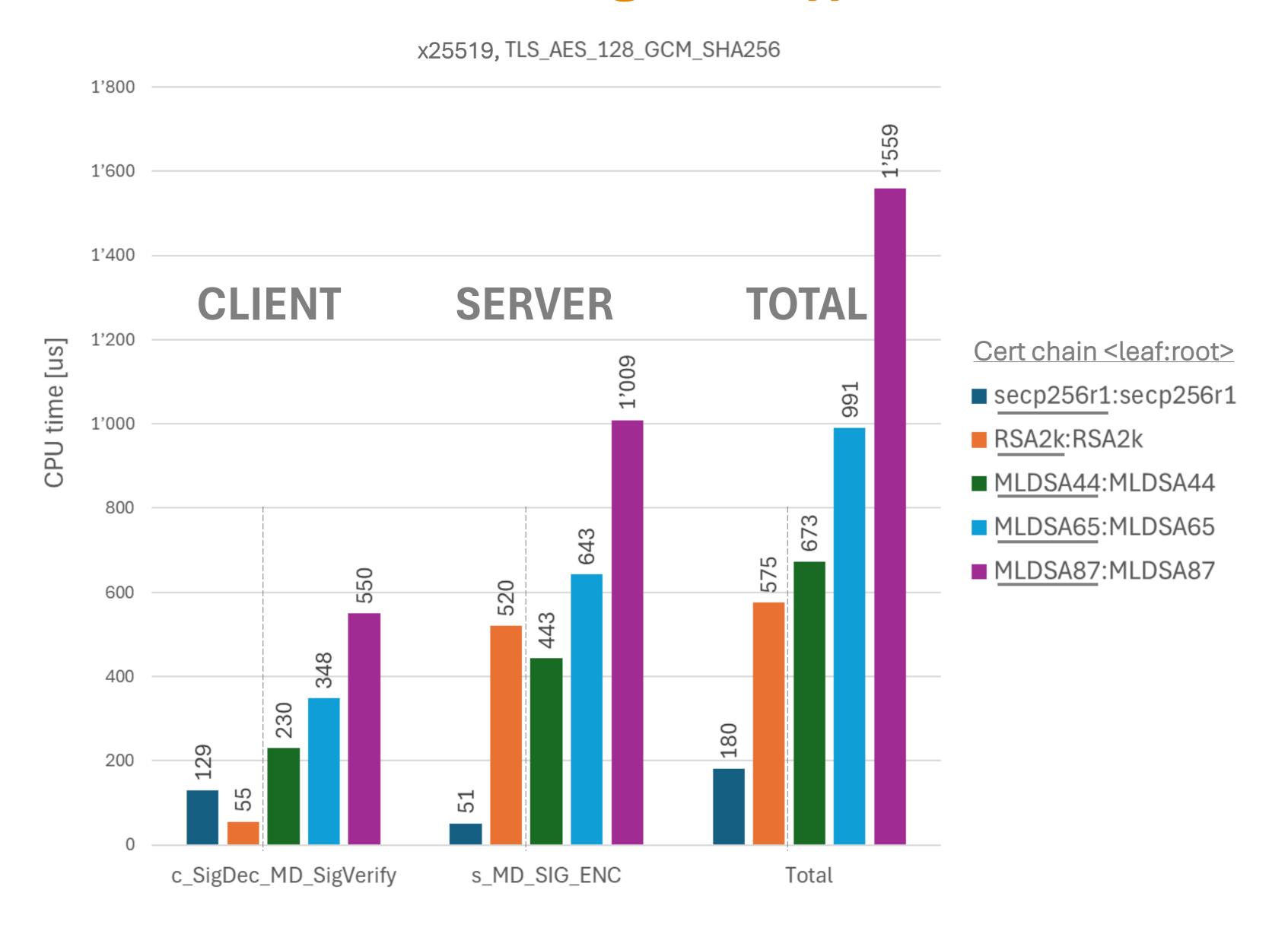


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Leaf certificates with RSA4k or Secp521r1 are very costly (up to +2 ms latency)

The root certificate algorithm has no influence on the timings; hence we zoom in on the next chart.....



## Measured Results: Sign/Verify Handshake Messages



#### Client:

- RSA2k is best
- MLDSA[44|65|87] adds  $\sim [200|300|400] \mu s$

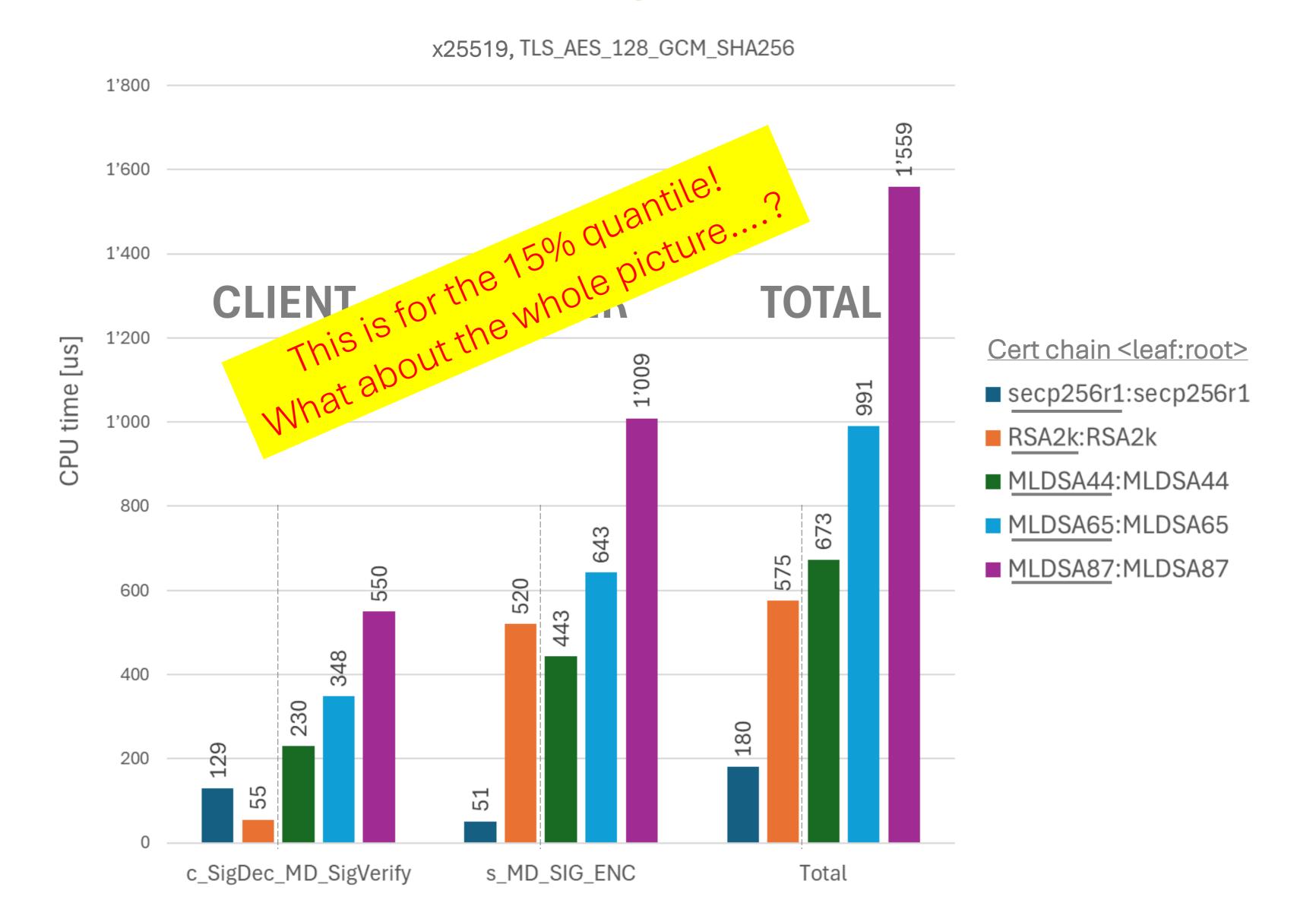
#### Server:

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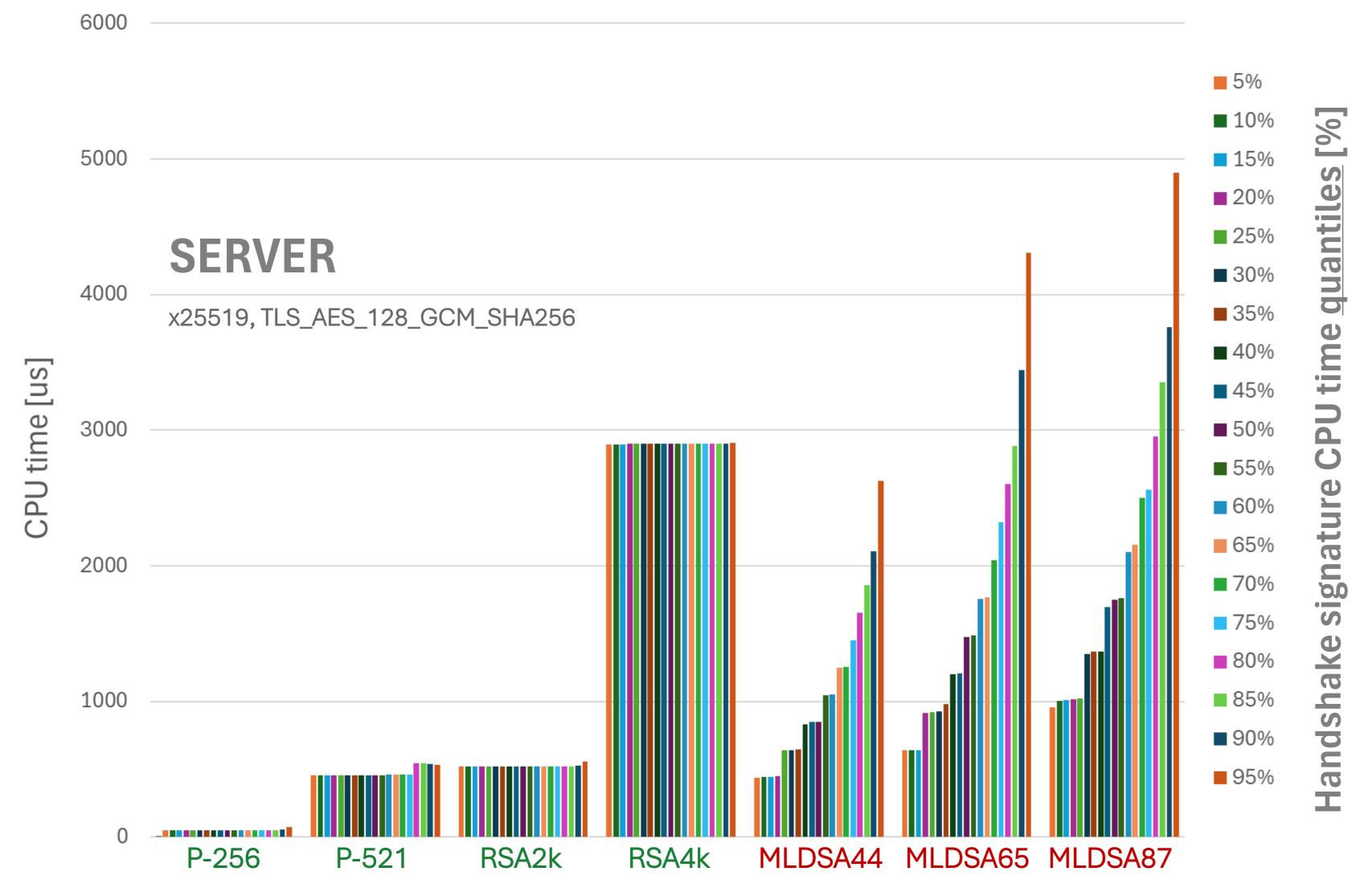
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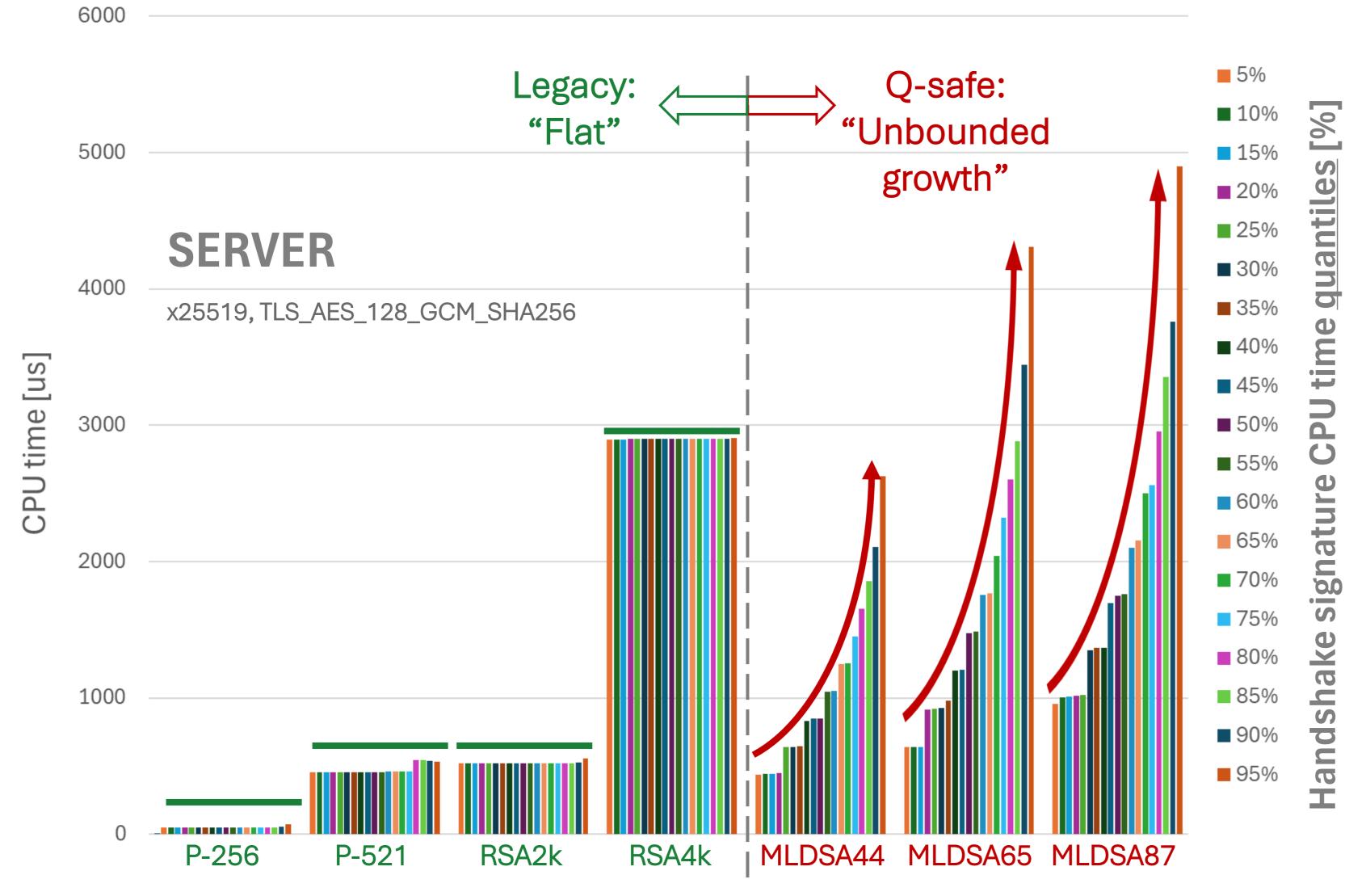
## Measured Results: <u>Sign</u> Handshake Messages



Significant tail latencies for MLDSA leaf certificates

Signature algorithm of the <u>leaf</u> certificate

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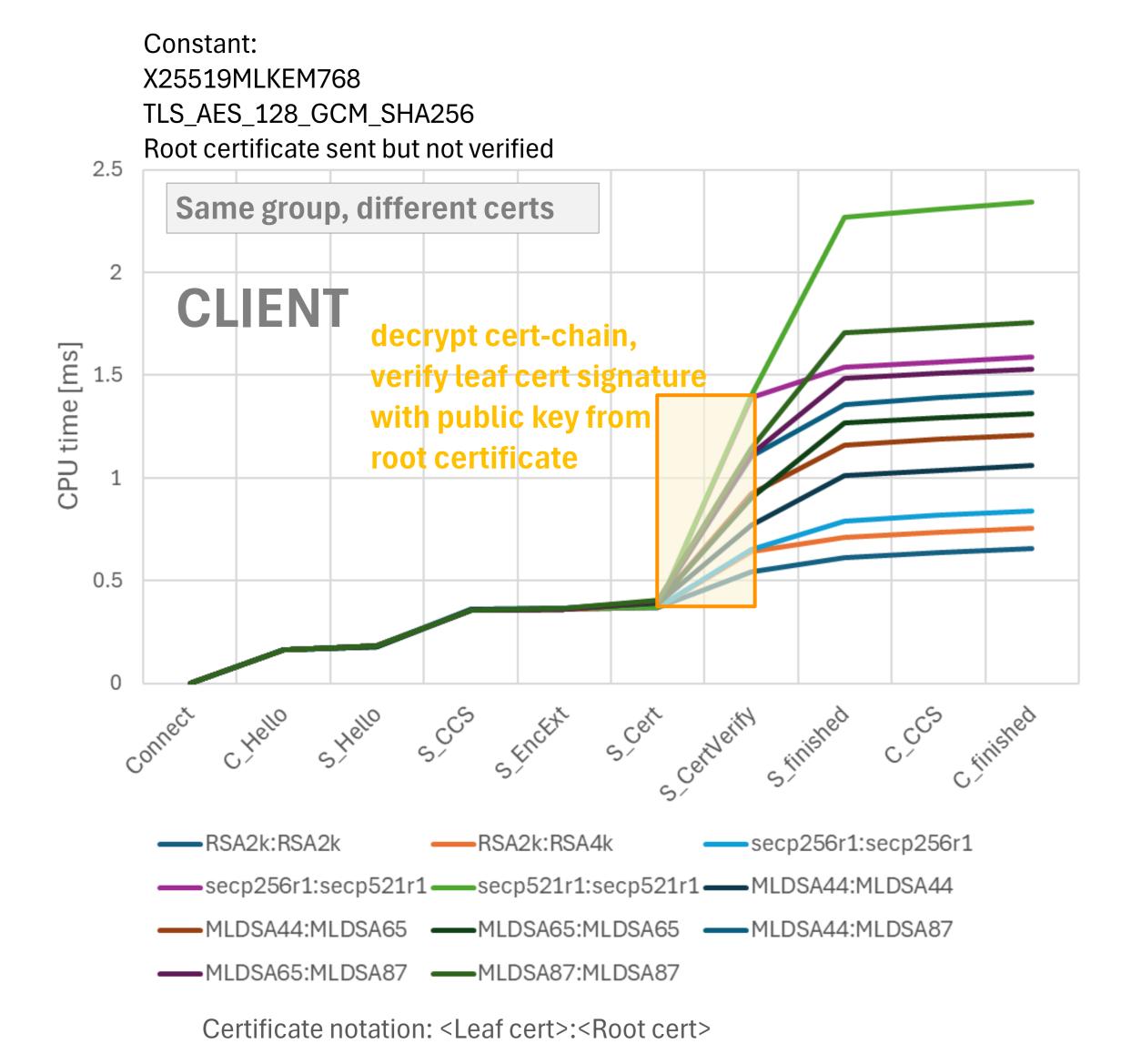
Significant tail latencies for MLDSA leaf certificates

This is happening on the server-side – and the server will ALWAYS sign.

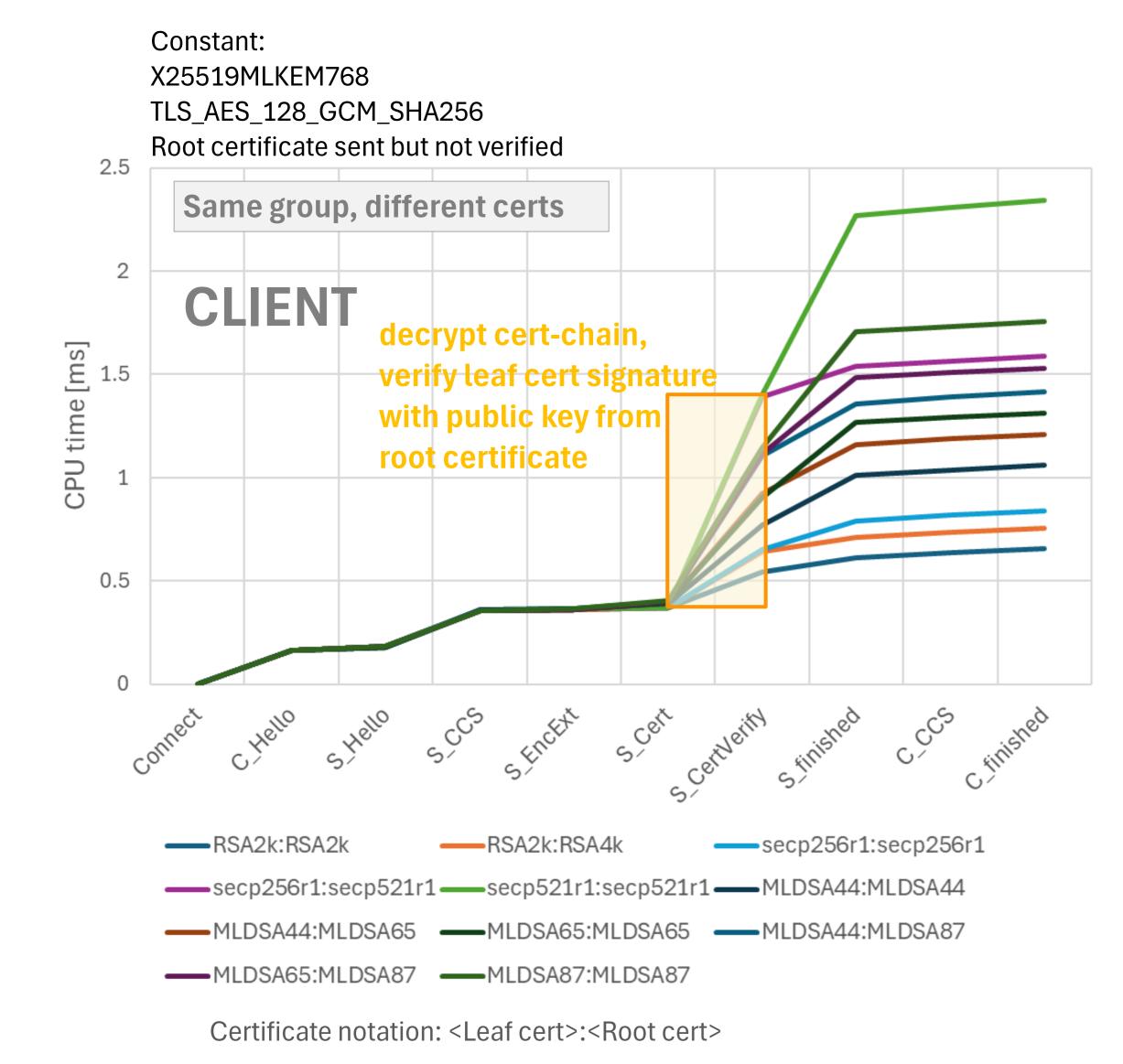
→ For a highly loaded server, ALL clients are affected by significant tail latencies

("highly loaded" = many parallel request, processed sequentially)

Signature algorithm of the <u>leaf</u> certificate



The leaf certificate signature algorithm is not involved, only the algorithm(s) of the parent cert(s) are used.



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#### Variations:

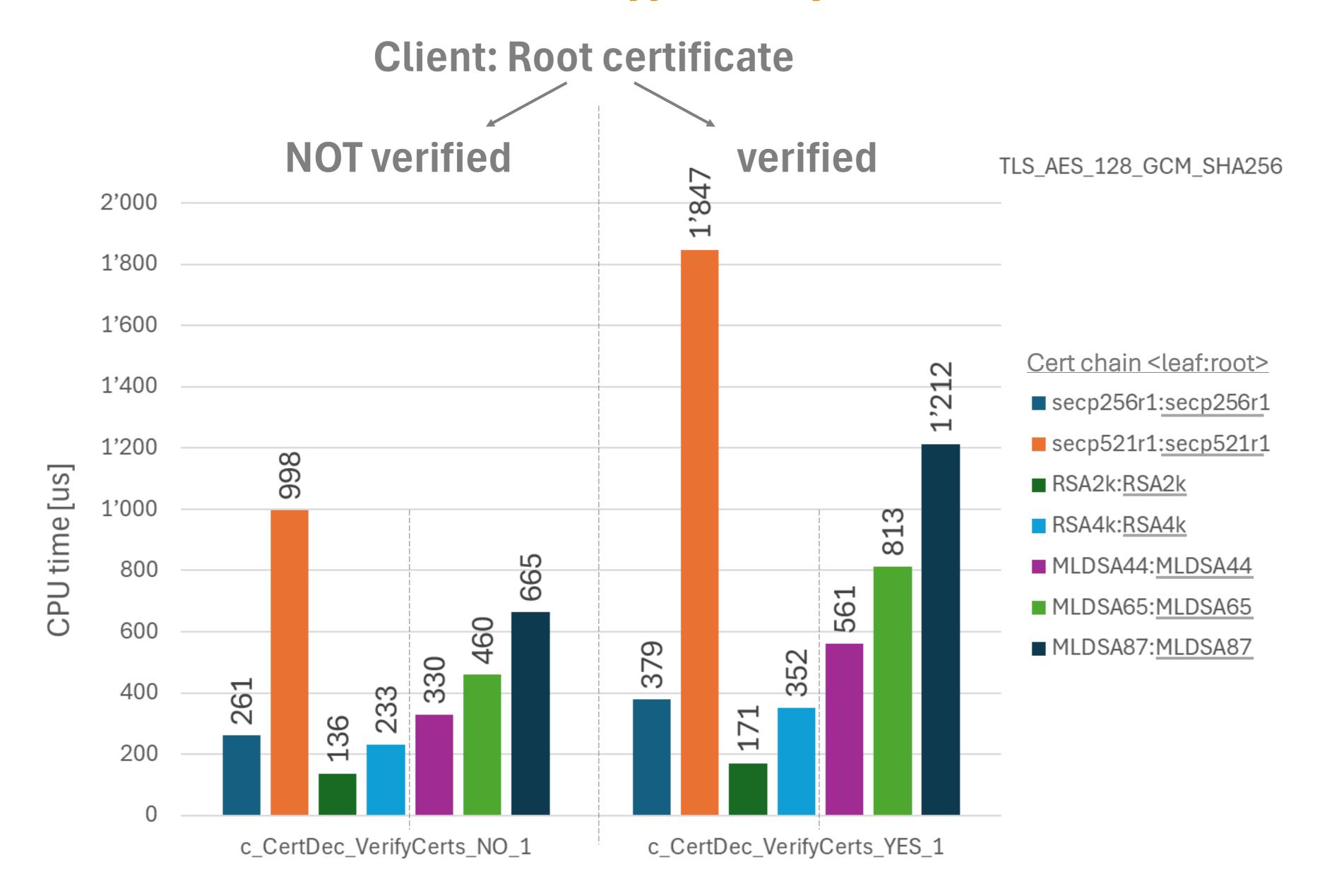
## The root certificate can optionally be sent

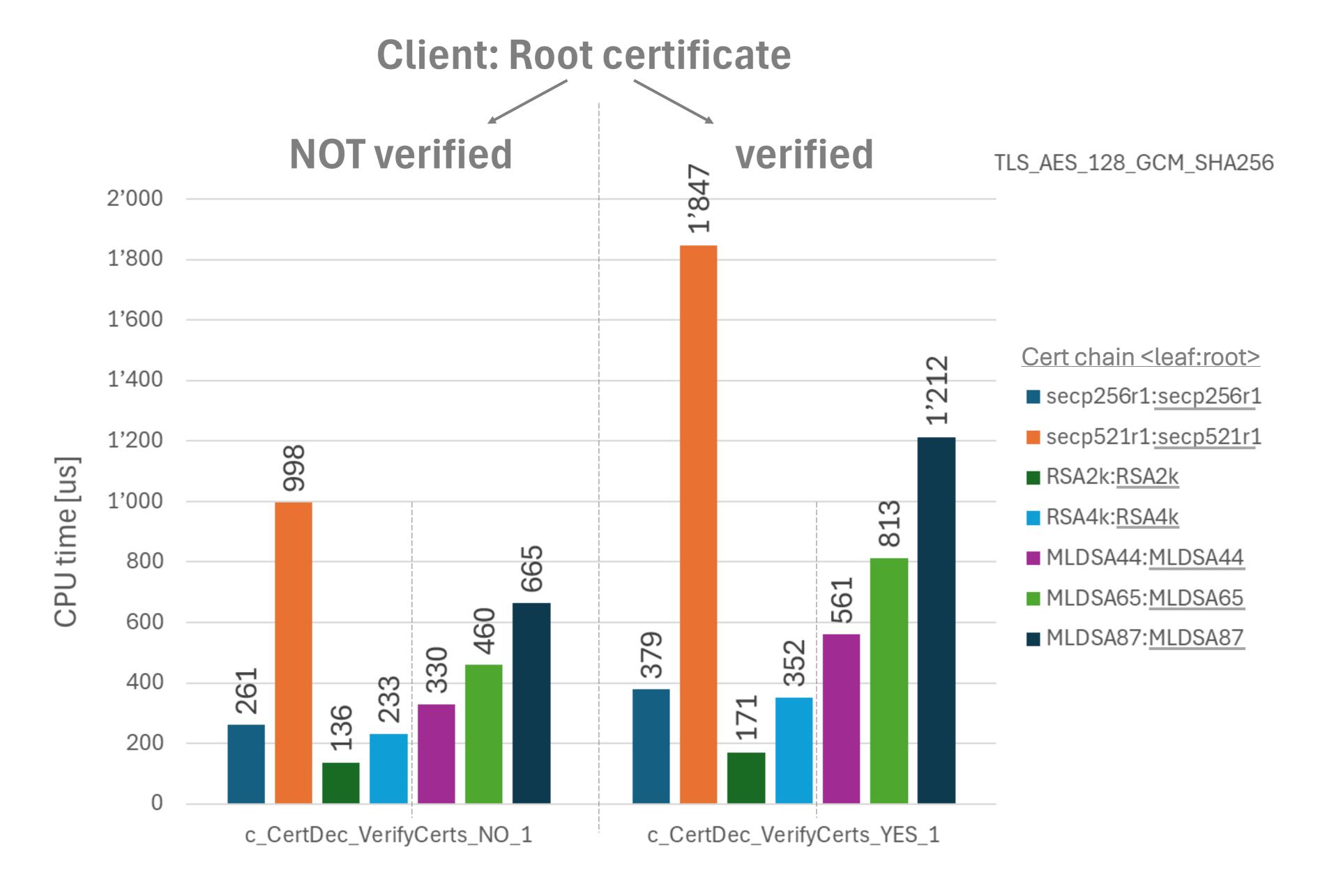
Latency penalties from receiving, storing and decrypting an extra certificate is  $\sim$ 45  $\mu$ s (± 15  $\mu$ s pending cipher), which is small in the overall scheme.

Therefore, we only show results for the case where the root certificate is not sent and when using cipher TLS\_AES\_128\_GCM\_SHA256.

## The root certificate can optionally be verified

This to capture the equivalent <u>case of an</u> intermediate certificate being in the chain.





All MLDSA versions beat a secp521r1 root parent certificate

Having to verify a parent (e.g. intermediate) certificate virtually doubles the latencies

Latencies for Secp256r1 vs RSA4k are virtually the same

Using MLDSA[44|65|87] instead of P-256|RSA4k adds  $\sim [100|200|400] \mu s$ 

## Overall Performance: Latencies

_			<b>15</b> %	% quant	ile	50% quantile			99% quantile		
ied	group	cert_chain	client	server	total	client	server	total	client	server	total
erif	x25519	secp256r1:secp256r1	0.644	0.263	0.906	0.653	0.265	0.917	0.705	0.356	1.061
<b>&gt;</b>	X25519MLKEM768	secp256r1:secp256r1	0.798	0.354	1.151	0.806	0.355	1.162	0.837	0.439	1.276
ON N	x25519	secp256r1:secp521r1	1.375	0.265	1.640	1.384	0.266	1.651	1.415	0.321	1.736
t	X25519MLKEM768	MLDSA65:MLDSA65	1.286	1.045	2.331	1.307	1.801	3.108	1.384	6.853	8.237
t Ce	X25519MLKEM768	MLDSA65:MLDSA87	1.446	1.072	2.518	1.464	1.807	3.271	1.563	6.945	8.508
ent	MLKEM1024	MLDSA87:MLDSA87	1.634	1.294	2.928	1.656	2.069	3.725	1.708	7.306	9.013
Par	SecP384r1MLKEM1024	MLDSA87:MLDSA87	2.505	2.143	4.648	2.524	2.920	5.444	2.871	8.135	11.006

### **OVERALL CPU TIMES** (aka LATENCY):

#### Groups & Certificates vs

- Quantiles,
- Parent cert verification

			15% quantile			50	% quan	tile	99% quantile		
	group	cert_chain	client	server	total	client	server	total	client	server	total
ed	x25519	secp256r1:secp256r1	0.764	0.265	1.029	0.781	0.266	1.047	0.825	0.357	1.182
rifi	X25519MLKEM768	secp256r1:secp256r1	0.905	0.353	1.258	0.918	0.355	1.274	0.973	0.375	1.348
ert ve	x25519	secp256r1:secp521r1	2.200	0.264	2.464	2.216	0.265	2.481	2.330	0.325	2.655
	X25519MLKEM768	MLDSA65:MLDSA65	1.635	1.057	2.693	1.657	1.808	3.465	1.840	6.932	8.772
ıtc	X25519MLKEM768	MLDSA65:MLDSA87	1.975	1.074	3.049	1.997	1.814	3.811	2.824	6.995	9.819
<u>re</u>	MLKEM1024	MLDSA87:MLDSA87	2.161	1.295	3.456	2.186	2.001	4.187	2.253	7.303	9.556
Pa	SecP384r1MLKEM1024	MLDSA87:MLDSA87	3.031	2.154	5.185	3.062	2.929	5.992	3.353	8.148	11.501

Only some selected group & certificate chain combinations shown

Cipher: TLS\_AES\_128\_GCM\_SHA256

Timings in [ms]

#### Rules of thumb for latencies:

- "Harvest now, decrypt later": Overall adds ~250  $\mu$ s (~=25%) latency, server load goes up ~33%
  - To put things into perspective: Having a secp521r1 certificate in the chain hurts ~3x more than using X25519MLKEM768
- Using MLDSA certificates really hurts performance:
  - $\sim$ 3x higher latency and server load in 50% quantile (compared to P-256),
  - up to ~10x higher latency and ~20x higher server load in 99% quantile

## Overall Performance: Data volume vs Bandwidths

			I	1	<u>Relevance:</u> new TLS session ≅ new client					1		
						Max sess	ion rate		]	Data transp	oort latency	/
			Data volume [E	vs data-center network BW [Msess/s]				not considering flight-time [us]				
Group	Cert chain	num certs	Client -> Server	Server -> Client	1 Gbps	10 Gbps	100 Gbps	400 Gbps	1 Gbps	10 Gbps	100 Gbps	400 Gbps
x25519	P-256:P-256	1	256	928	0.11	1.08	10.78	43.10	11.84	1.18	0.12	0.03
MLKEM512	P-256:P-256	1	1016	1665	0.06	0.60	6.01	24.02	26.81	2.68	0.27	0.07
MLKEM768	P-256:P-256	1	1400	1985	0.05	0.50	5.04	20.15	33.85	3.39	0.34	0.08
X25519MLKEM768	P-256:P-256	1	1432	2017	0.05	0.50	4.96	19.83	34.49	3.45	0.34	0.09
MLKEM1024	P-256:P-256	1	1784	2464	0.04	0.41	4.06	16.23	42.48	4.25	0.42	0.11
MLKEM512	MLDSA44:MLDSA44	1	1014	7612	0.01	0.13	1.31	5.25	86.26	8.63	0.86	0.22
MLKEM768	MLDSA65:MLDSA65	1	1398	10350	0.01	0.10	0.97	3.86	117.48	11.75	1.17	0.29
X25519MLKEM768	MLDSA65:MLDSA65	1	1430	10382	0.01	0.10	0.96	3.85	118.12	11.81	1.18	0.30
MLKEM1024	MLDSA87:MLDSA87	1	1782	14106	0.01	0.07	0.71	2.84	158.88	15.89	1.59	0.40
MLKEM512	MLDSA44:MLDSA65	2	1016	14060	0.01	0.07	0.71	2.84	150.76	15.08	1.51	0.38
MLKEM768	MLDSA65:MLDSA87	2	1400	19185	0.01	0.05	0.52	2.08	205.85	20.59	2.06	0.51
X25519MLKEM768	MLDSA65:MLDSA87	2	1432	19217	0.01	0.05	0.52	2.08	206.49	20.65	2.06	0.52
MLKEM1024	MLDSA87:MLDSA87	2	1782	21623	0.00	0.05	0.46	1.85	234.05	23.41	2.34	0.59

Latencies due to data transport latencies are relatively small and hence of no big concern: max +23 μs @ 10Gbps Data transfers Server -> Client are dominating

The data-center egress network BW becomes the limiting factor for TLS session establishment rates Rules of thumb for session establishment rates:

- Protecting against 'harvest now, decrypt later' lowers upper bound for session establishment capability by factor of ~2
- Adding a **2-stage Q-safe cert chain**, without sending the root certificate, cuts it by ~10x
- When also sending the parent e.g. intermediate certificate, the cut is ~20x

## Insights Summary

#### **CIPHERS**

Choice of ciphers has a small impact on overall latency (variations are within  $\pm$  50  $\mu$ s)

#### **GROUPS**

Protecting against 'harvest now, decrypt later'

- "Pure" Q-safe key agreement algorithms are virtually at par with their legacy counterparts
- MLKEM512 (128 sec bits) is even faster than legacy groups
- Hybrid X25519MLKEM768 adds ~250 μs (~=25%) to overall latency, Server load goes up ~33%

### CERTIFICATES

Using Q-safe algorithms for authentication

- MLDSA is costly per-se AND leads to unbound tail latencies:
  - ~3x higher latency in 50% quantile,
  - BUT up to  $\sim$ 10x higher latency and  $\sim$ 20x higher server load in 99% quantile
- Having to verify a parent (e.g. intermediate) certificate virtually doubles the latencies  $\rightarrow$  use 2-stage cert chains and don't send the root certificate

#### DATA VOLUMES

Latencies due to data transport latencies are relatively small and hence of no big concern: max +23 µs @ 10Gbps The data-center egress network BW becomes the limiting factor for TLS session establishment rates Rules of thumb:

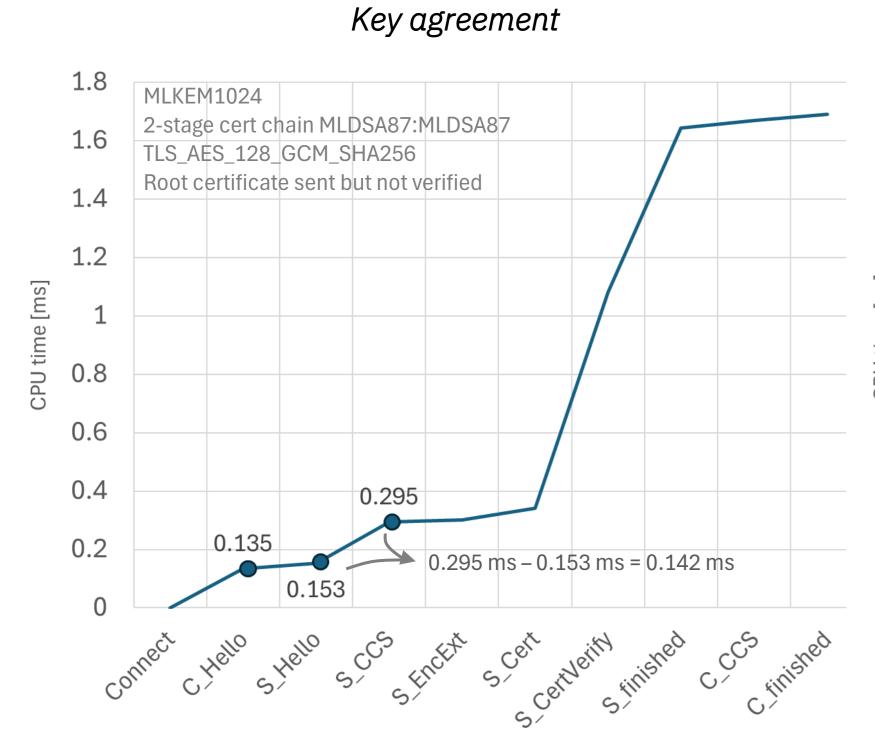
- Protecting against 'harvest now, decrypt later' is cutting session establishment capability by factor of ~2
- Adding a 2-stage Q-safe cert chain, without sending the root certificate, cuts session rates by  $\sim 10x$
- When also sending the parent e.g. intermediate certificate, the cut is ~20x

# THANK YOU!

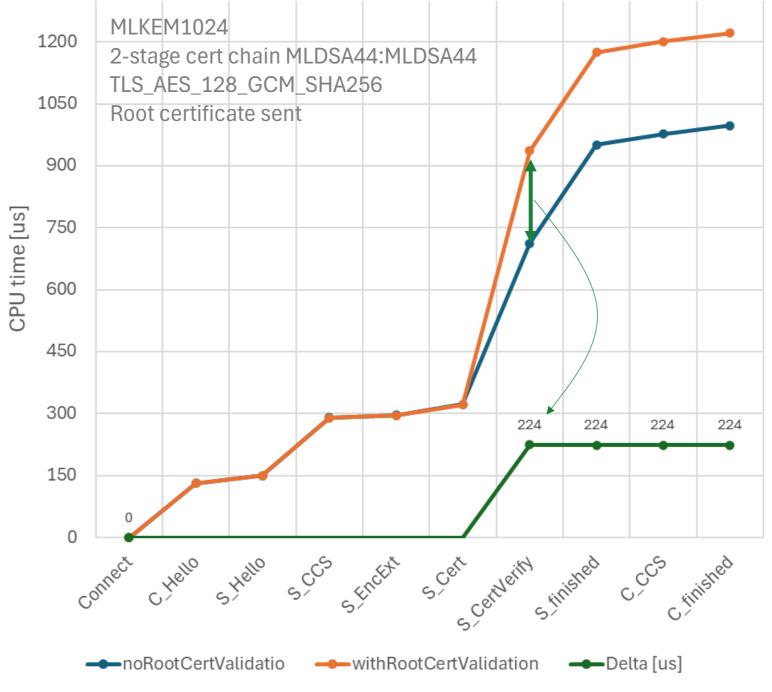
# Backup

## Data extraction: Examples (client-side)

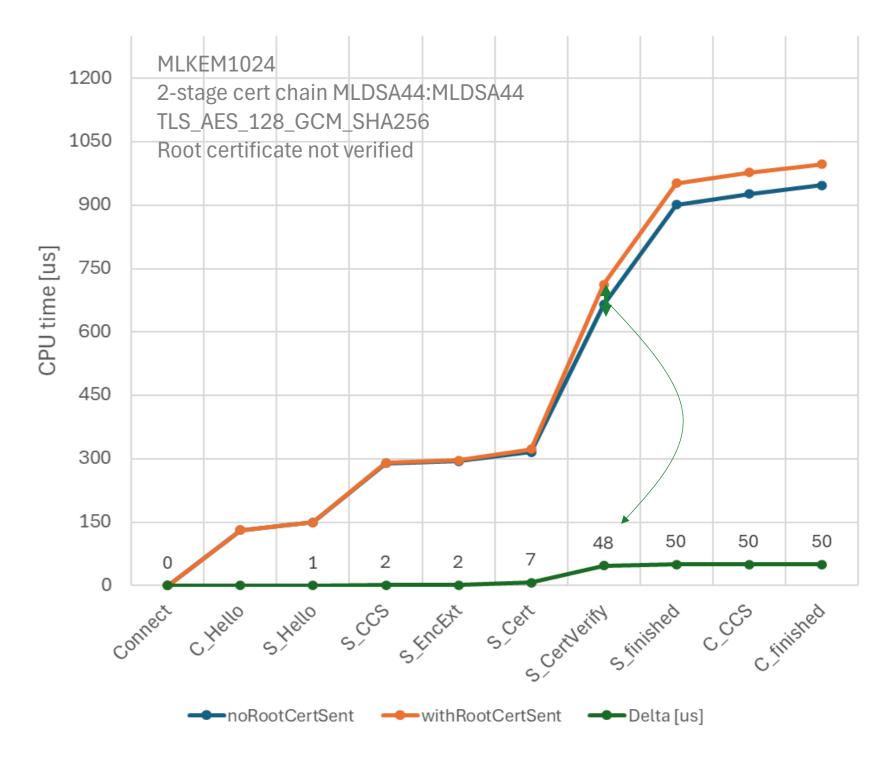
Data extraction can be based on a single measurement OR based on the difference between two measurements







#### Impact of sending the root certificate



Insight: It takes 135 µs for the client to generate a MLKEM1024 key pair and 142 μs to derive a common MLKEM1024 secret

Insight: It takes 224 µs for the client to verify the (self-)signature of a MLDSA44 root certificate

Insight: It takes 48 µs for the client to decrypt a MLDSA44 root certificate plus do the bookkeeping

Constant: CLIENT, MLKEM1024, TLS\_AES\_128\_GCM\_SHA256, MLDSA44:MLDSA44

Note: Generated with OpenSSL v3.5.0

## Crypto-Algorithms: Available Options



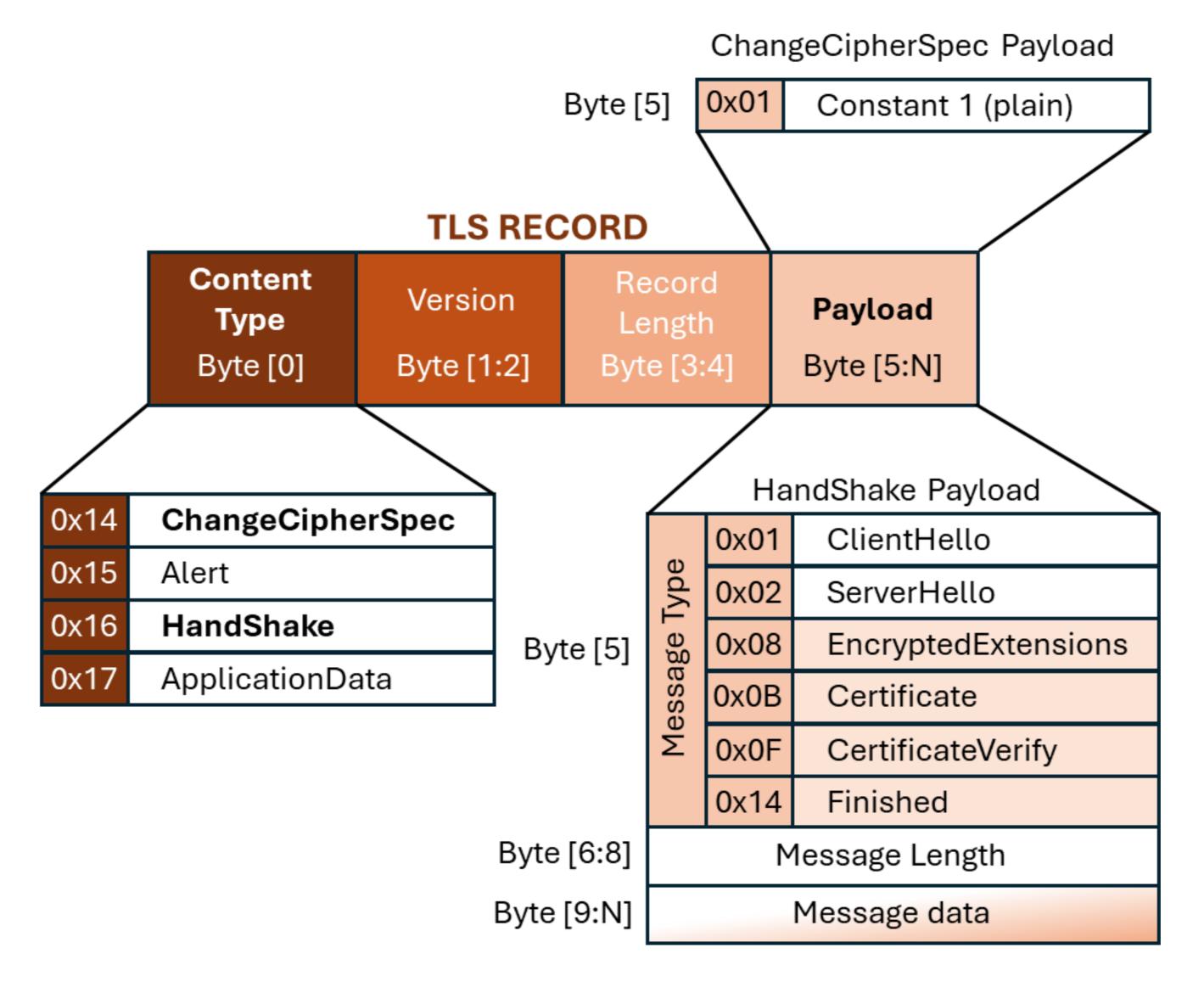
	Groups	Certificate algorithms [leaf:root]	Ciphersuites	Root certificate sent	Root certificate verified	
	MLKEM512	RSA2k:RSA2k	TLS_AES_128_GCM_SHA256	YES	YES	
	MLKEM768	RSA2k:RSA4k	TLS_AES_256_GCM_SHA384	NO	NO	
	MLKEM1024	secp256r1:secp256r1	TLS_CHACHA20_POLY1305_SHA256	П	П	
	x25519	secp256r1:secp521r1	TLS_AES_128_CCM_SHA256			
	secp256r1	secp521r1:secp521r1	TLS_AES_128_CCM_8_SHA256_			
	X25519MLKEM768	MLDSA44:MLDSA44				
	SecP256r1MLKEM768	MLDSA44:MLDSA65			Derive dura	ation to verify the
	x448	MLDSA65:MLDSA65			signature o	n a certificate
	secp384r1	MLDSA44:MLDSA87				
	/> secp521r1	MLDSA65:MLDSA87		7		
1		MLDSA87:MLDSA87 <		Deriv	e duration to	encrypt or
<u>ا</u>				decry	pt a certifica	ate

Derive duration for

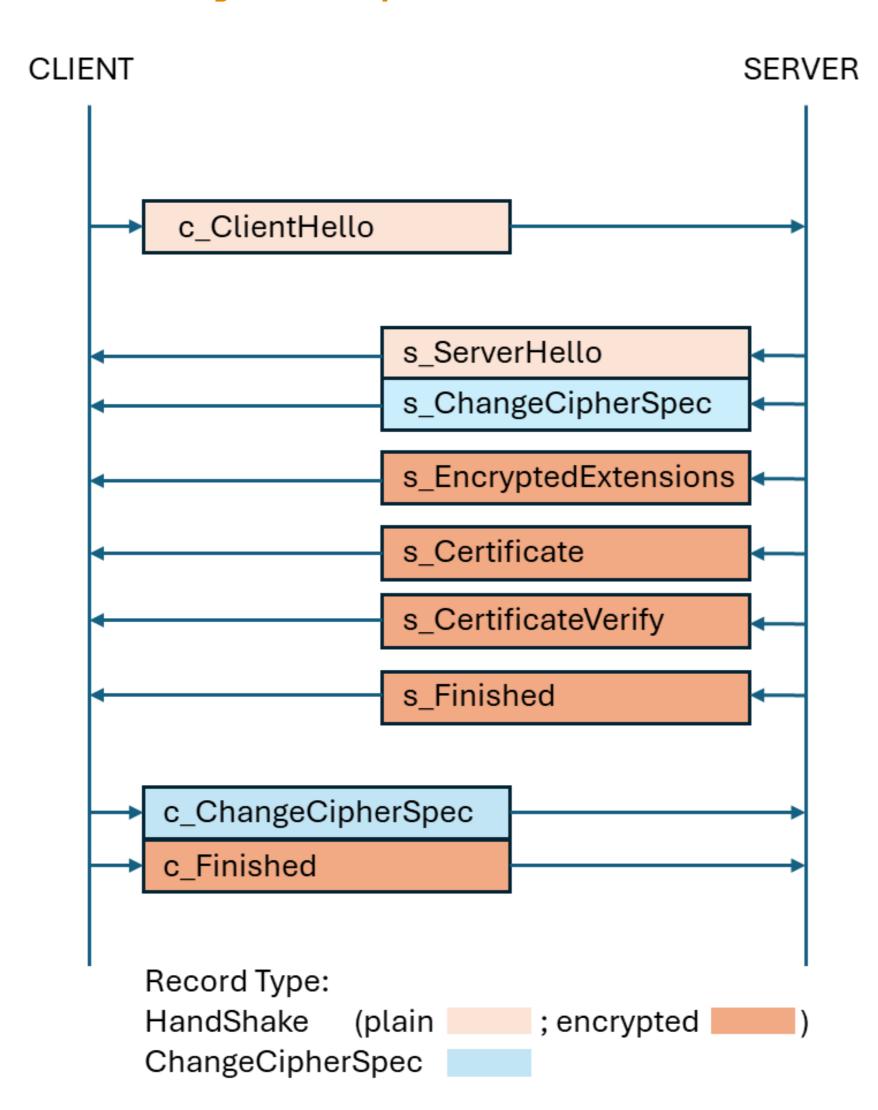
- generating key-pair(s)
- generating common secret

Derive duration to create a handshake digest and to sign or verify a signature

## Anatomy of a TLS v1.3 Record



# TLS v1.3 session establishment sequence without HelloRetryRequest (HRR)



#### Note:

No HRR implies that the server can immediately accept a key-share offered by the client